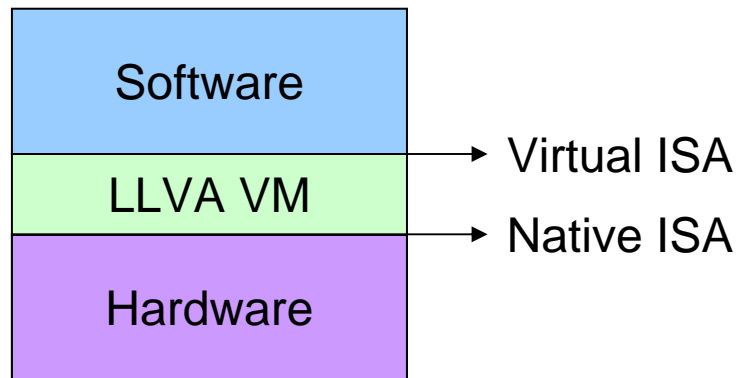


# Secure Virtual Architecture

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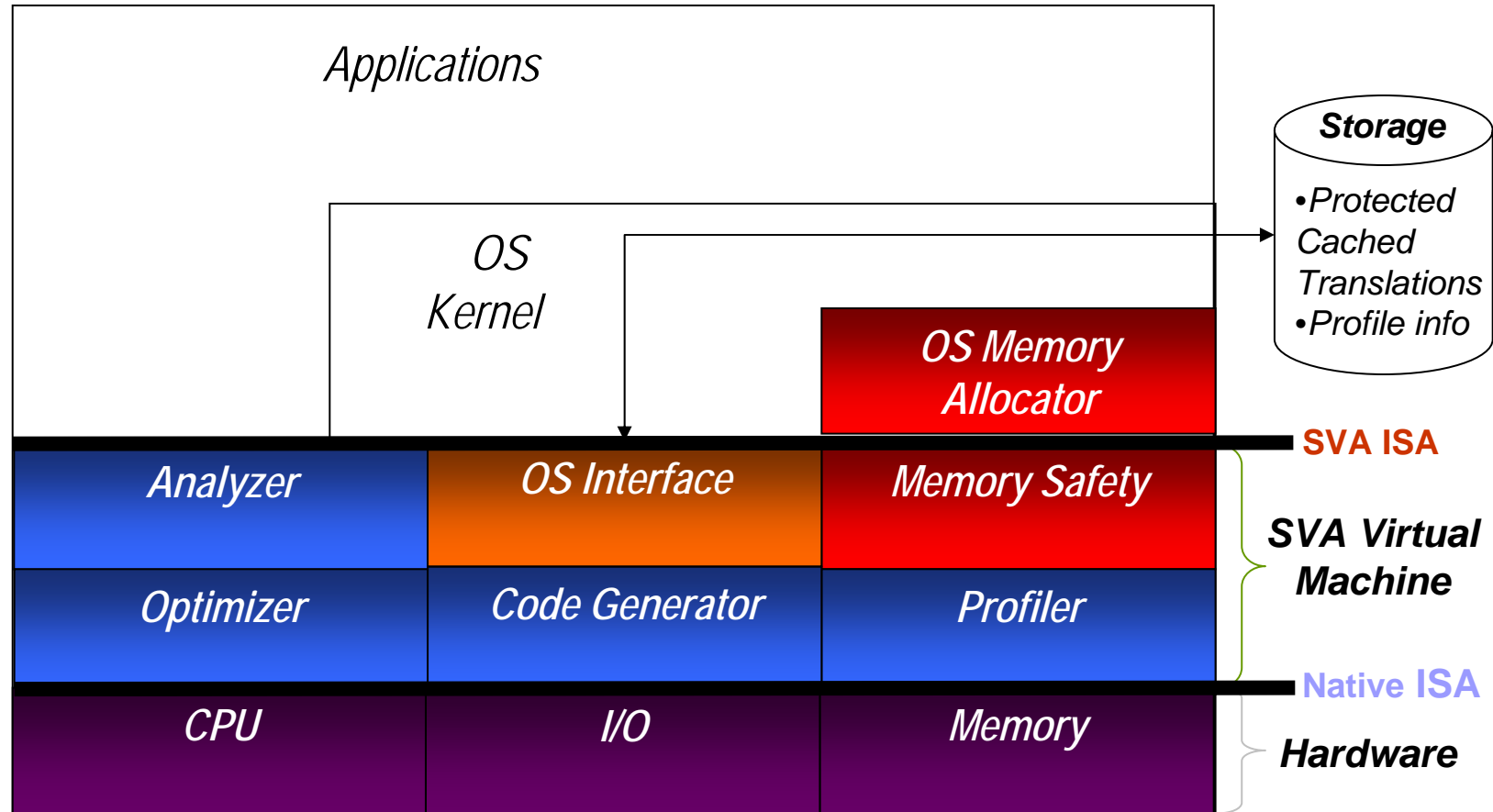


# Secure Virtual Machine



- What is it?
  - A *compiler-based* virtual machine running *below* an operating system
  - Enforces security policies for *all* software (including the OS kernel)
- Why build it?
  - Kernel code can be exploited
  - Allows analysis across traditional boundaries

# SVA System Architecture



# SVA: OS Interface

- **Kernels require new functionality**
  - **Hardware Control**
    - Performing I/O
    - Installing interrupt handlers
  - **State Manipulation**
    - Context switching
    - Signal handler dispatch

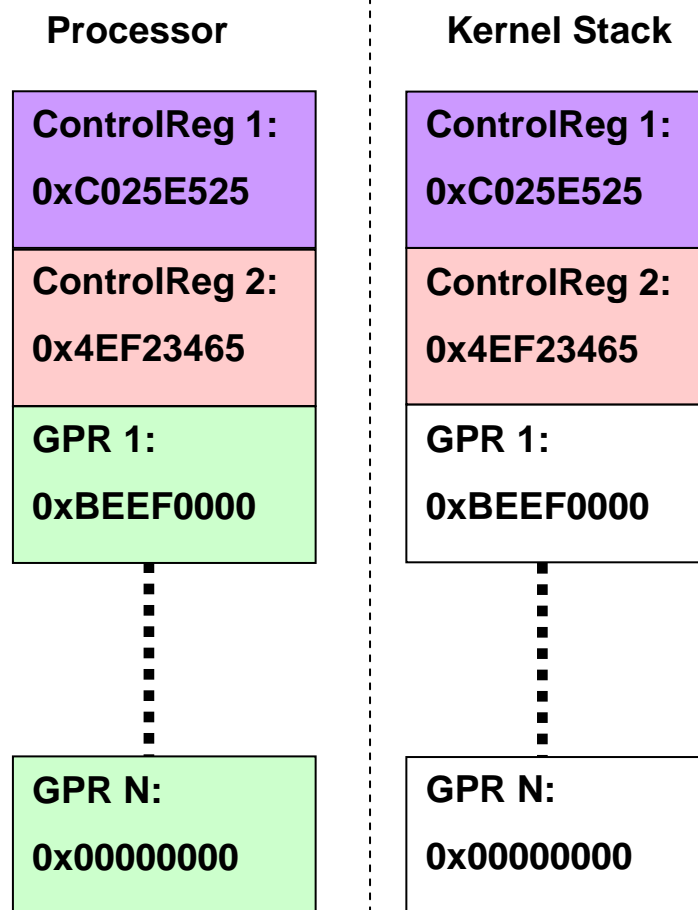
# Hardware Control

- Registration functions
  - void *llva\_register\_syscall* (int number, int (\*f)(void \* icontext, ...))
  - void *llva\_register\_interrupt* (int number, int (\*f)(void \* icontext))
  - void *llva\_register\_exception* (int number, int (\*f)(void \* icontext))
- I/O
  - int *llva\_io\_read* (ioptr\_t ioaddress)
  - void *llva\_io\_write* (ioptr\_t ioaddress, int value)
- Atomic Operations
  - int *llva\_swap\_and\_phi* (void \* address, int value)
  - int *llva\_compare\_and\_swap* (void \* address, int compare, int value)
- Memory Management
  - void *llva\_load\_pgtable* (void \* table)
  - void \* *llva\_save\_pgtable* ()

# State Manipulation

- Allow OS to see the *existence* of native state
- OS does not understand the semantics of native state

# Lazy State Saving on Interrupt



- How to take advantage of low latency interrupt facilities?
  - shadow registers (e.g. ARM)
  - register windows (e.g. SPARC)
- On interrupt, SVM saves subset of processor state on the kernel stack
- Can leave state in registers if kernel does not overwrite it
- Kernel can commit all state to memory if required
- Pointer to Interrupt Context passed to system call, interrupt, and trap handlers

# Manipulating Interrupt Context

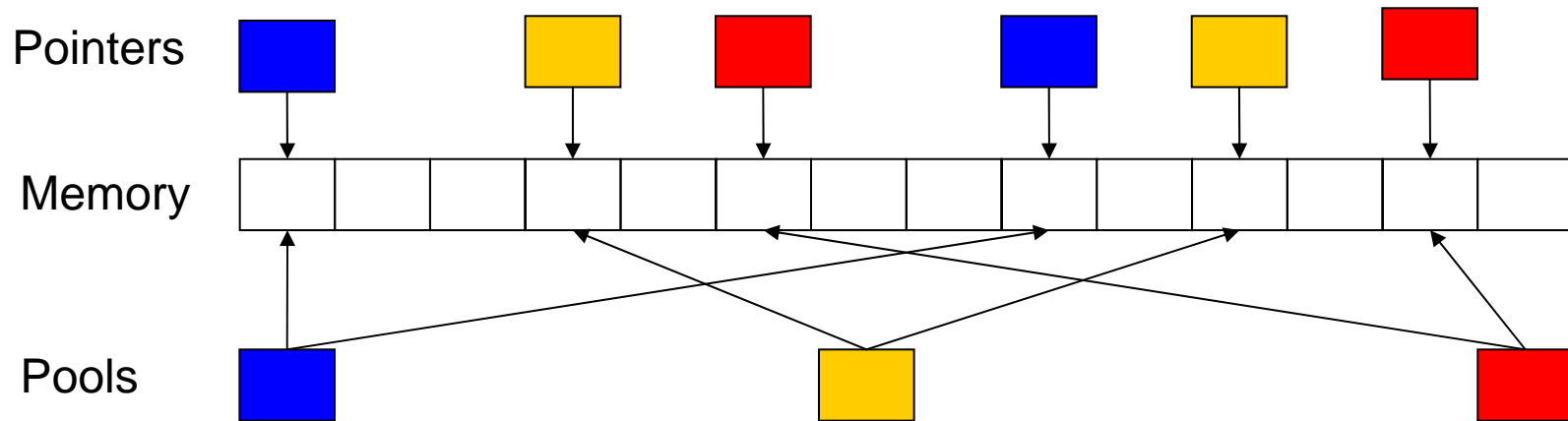
- Interrupt Context  $\leftrightarrow$  Memory
  - void *llva\_icontext\_save* (void \* icontext, void \* buffer)
  - void *llva\_icontext\_load* (void \* icontext, void \* buffer)
- Commit
  - void *llva\_icontext\_commit* (void \* icontext)
- Push function frames
  - void *llva\_ipush\_function* (void \* icontext, void (\*f)(...), ...)



# Manipulating Processor State

- Context Switching (manipulates current state)
  - void *llva\_save\_integer* (void \* buffer)
  - void *llva\_load\_integer* (void \* buffer)
  - void *llva\_save\_fp* (void \* buffer, bool save\_always)
  - void *llva\_load\_fp* (void \* buffer)

# SVA: Memory Safety for OS Kernels



- Use static analysis to prove safe memory accesses
- Use alias analysis (DSA) to group objects into logical pools
- Virtual machine records object allocations in pools
- Run-time checks only check objects in a single pool

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  - Rob Bocchino
  - Brent Monroe
  - Pierre Salverda
  - David Raila
- Vikram Adve
- The LLVM Community

# References (<http://llvm.org/pubs>)

- SVA
  - MICRO '03
  - VEE '06
  - WIOSCA '06
- Data Structure Analysis (DSA)
  - PLDI '07
- Automatic Pool Allocation
  - MSP '02
  - PLDI '05
  - Lattner PhD Thesis
- Memory Safety
  - LCTES '03
  - TECS '05
  - CASES '02
  - Dhurjati PhD Thesis

Questions?

