

Polly

Polyhedral Transformations for LLVM

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Outline

1 The Polyhedral Model

2 Research Projects

3 Polly

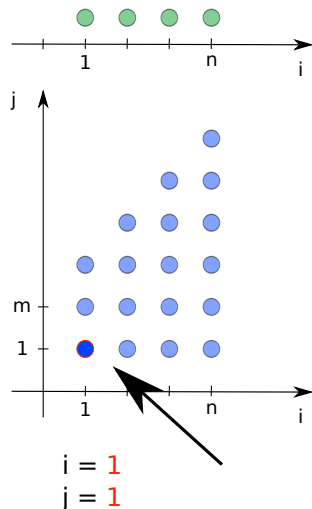
The SCoP - Static Control Part

```
for i = 1 to (5n + 3)
  for j = n to (4i + 3n + 4)
    A[i-j] = A[i]
  if i < (n - 20)
    A[i+20] = j
```

- Structured control flow
 - ▶ Regular for loops
 - ▶ Conditions
- Affine expressions in induction variables and parameters for:
 - ▶ Loop bounds, conditions, access functions
- Side effect free

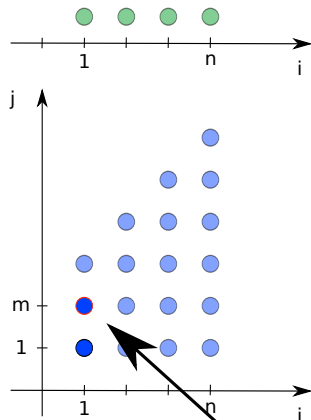
The Iteration Domain / Schedule

```
for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```



The Iteration Domain / Schedule

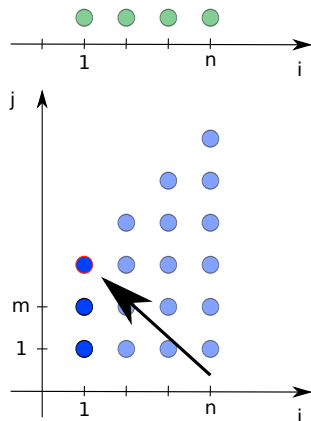
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for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
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}
```



$i = 1$
 $j = 2 = m$

The Iteration Domain / Schedule

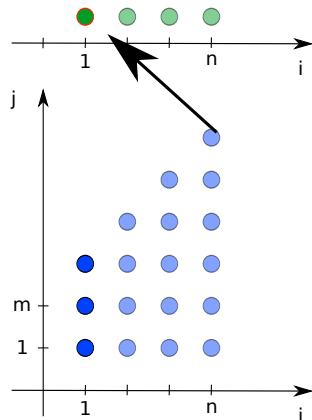
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}
```



$$i = 1$$
$$j = 3 = (i + m)$$

The Iteration Domain / Schedule

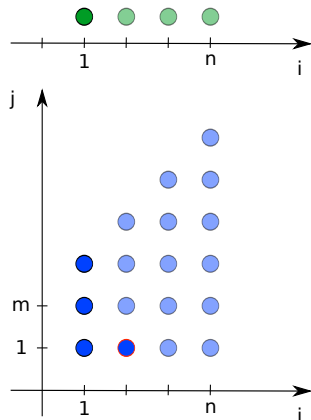
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}
```



$i = 1$
 $j = n/a$

The Iteration Domain / Schedule

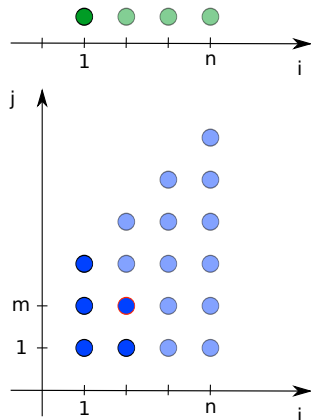
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}
```



$i = 2$
 $j = 1$

The Iteration Domain / Schedule

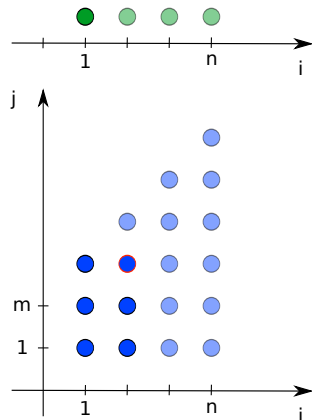
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for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```



$i = 2$
 $j = 2$

The Iteration Domain / Schedule

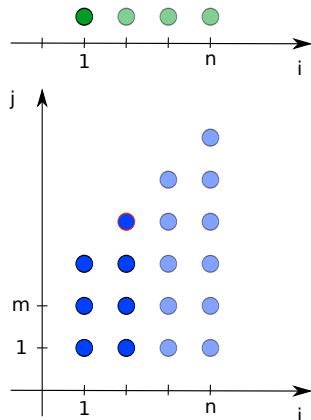
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}
```



$i = 2$
 $j = 3$

The Iteration Domain / Schedule

```
for (i = 1; i <= n; i++) {  
  for (j = 1; j <= i + m; j++)  
    A[i][j] = A[i-1][j] + A[i][j-1]  
  
  A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```

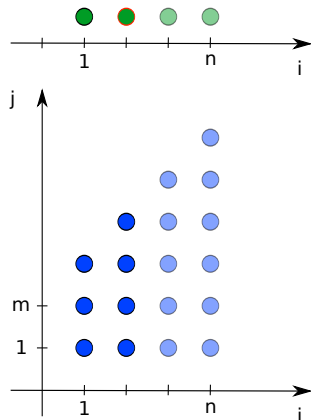


$$i = 2$$

$$j = 4 = (i + m)$$

The Iteration Domain / Schedule

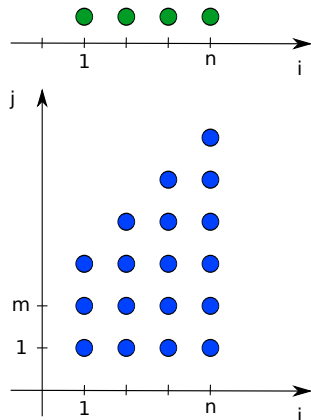
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    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
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}
```



$i = 2$
 $j = n/a$

The Iteration Domain / Schedule

```
for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```



The Iteration Domain / Schedule

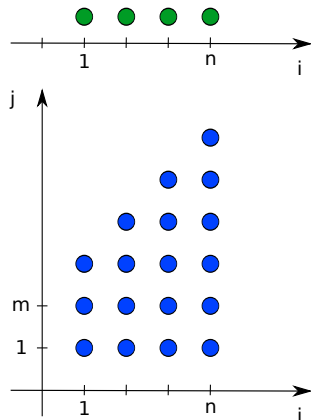
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    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```

Domains

$$1 \leq i \leq n$$

$$1 \leq j \leq i + m$$

$$1 \leq i \leq n$$



The Iteration Domain / Schedule

```
for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
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}
```

Domains

$$1 \leq i \leq n$$

$$1 \leq j \leq i + m$$

$$1 \leq i \leq n$$

Schedules

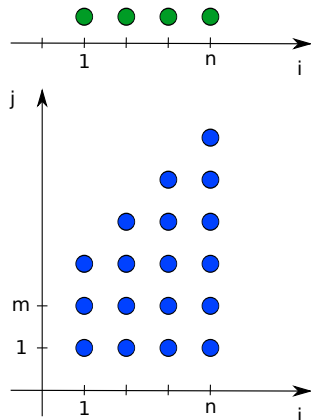
$$s_1 = i$$

$$s_2 = 0$$

$$s_3 = j$$

$$s_1 = i$$

$$s_2 = 1$$



The Iteration Domain / Schedule

```
for (i = 1; i <= n; i++) {  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```

Domains

$$1 \leq i \leq n$$

$$1 \leq j \leq i + m$$

$$1 \leq i \leq n$$

Schedules

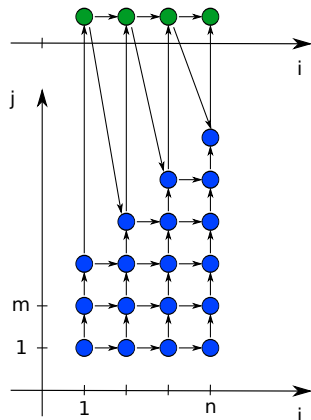
$$s_1 = i$$

$$s_2 = 0$$

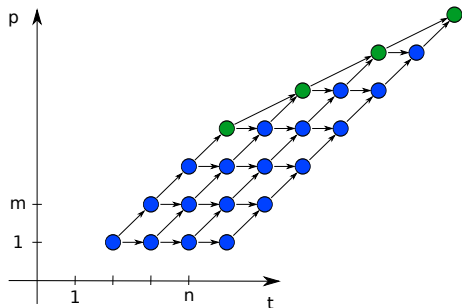
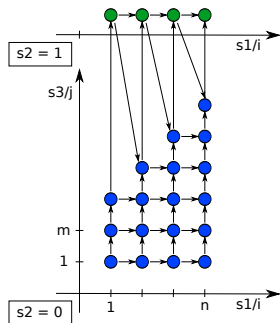
$$s_3 = j$$

$$s_1 = i$$

$$s_2 = 1$$



Scheduling Transformation



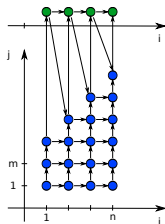
Original Schedules

$$\begin{array}{ll} s_1 = i & s_1 = i \\ s_2 = 0 & s_2 = 1 \\ s_3 = j & \end{array}$$

Transformed Schedules

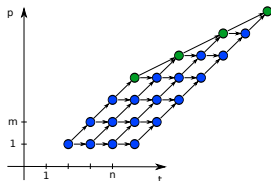
$$\begin{array}{ll} p = j & p = i + m + 1 \\ t = i + j & t = 2i + m + 1 \end{array}$$

Code Generation



Original Code

```
for (i = 1; i <= n; i++) {  
  
    for (j = 1; j <= i + m; j++)  
        A[i][j] = A[i-1][j] + A[i][j-1]  
  
    A[i][i+m+1] = A[i-1][i+m] + A[i][i+m]  
}
```



Transformed Code

```
parfor (p = 1; p <= m+n+1; p++) {  
    if (p >= m+2)  
        A[p-m-1][p] = A[p-m-2][p-1]  
    for (t = max(p+1, 2*p-m); t <= p+n; t++)  
        A[-p+t][p] = A[-p+t-1][p] + A[-p+t][p-1]  
}
```

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Source to source frameworks

Suif | Omega | LooPo | Pluto | ...

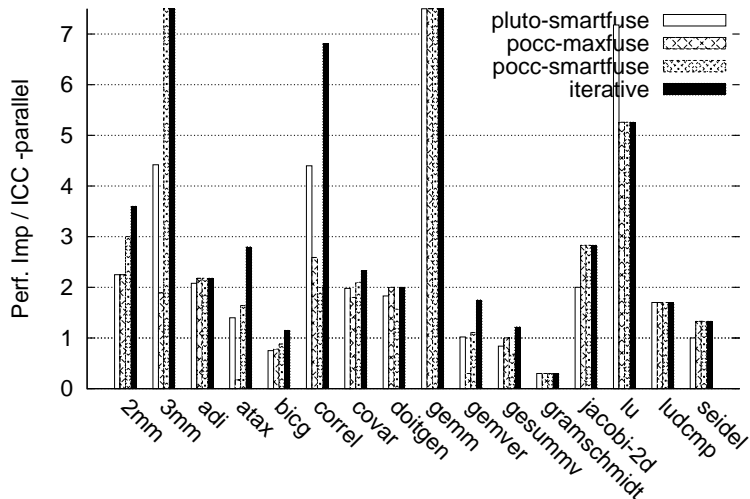
- More than 20 years of research
- Work on
 - ▶ Tiling / Parallelization / Prevectorization
 - ▶ CUDA / GPGPU
 - ▶ Coarse grain parallelism / Grid computing
- Advanced algorithms / Proven performance

Speed comparison

Compile mode	GCC 4.5.1	clang 2.8	ICC 11.1
-O3	1m 22.0s	1m 22.0s	0m 5.8s
pluto-tiled -O3	0m 6.1s	0m 5.8s	0m 2.5s

Table: Matrix multiplication on Intel i5 M 520 - *double* 2048 x 2048

Performance Improvement - AMD Opteron 8380 (16 threads)



Source to source frameworks

Restrictions

- Limited to subset of C/C++
- Require annotated C code
- Only canonical code
- Correct? (Integer overflow, Operator overloading, ...)

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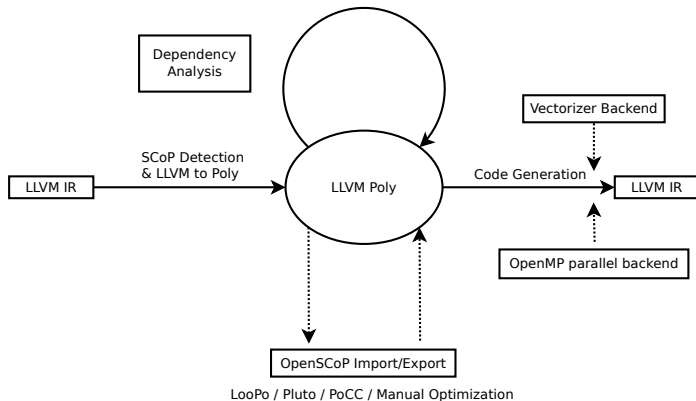
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The Architecture

Transformations

- * Classical loop transformations (Blocking, Interchange, Fusion, ...)
- * Expose parallelism
- * Dead instruction elimination / Constant propagation



Why optimize on LLVM-IR?

- Frontend independent
- Fully automatic
- High SCoP coverage
- Integration with vectorizer/ OpenMP code generator

The SCoP - Revised

Thanks to

- Scalar evolution
- Loop/Region detection
- LLVM canonicalization passes

SCoP - The LLVM way

- ~~Limited Scalars~~ Any scalars
- Structured control flow
 - ▶ ~~Regular for loops~~ Anything that acts like a regular for loop
 - ▶ Conditions
- ~~Affine expressions~~ Expressions that calculate an affine result
- Side effect free known
- ~~Memory accesses only through arrays~~

Valid SCoPs

do..while loop

```
i = 0;

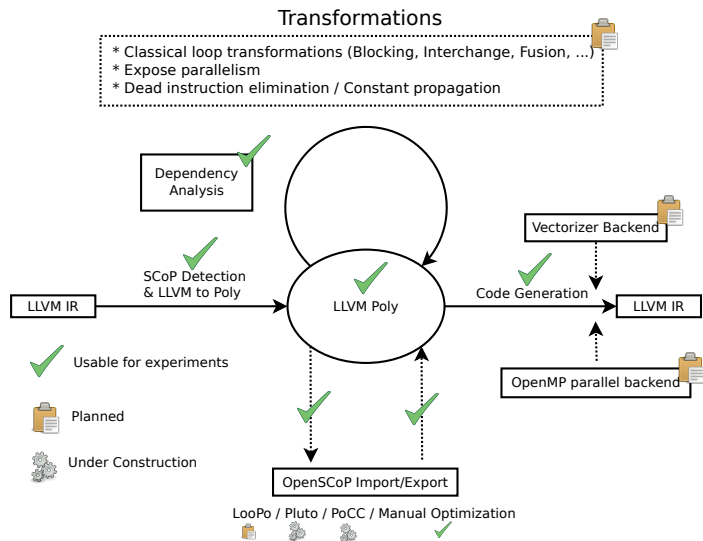
do {
    int b = 2 * i;
    int c = b * 3 + 5 * i;
    A[c] = i;
    i += 2;
} while (i < N);
```

pointer loop

```
int A[1024];

void pointer_loop () {
    int *B = A;
    while (B < &A[1024]) {
        *B = 1;
        ++B;
    }
}
```

The Architecture



Open Topics



- Multi dimensional arrays - Delinearization
- Integer modulo arithmetics
- Optimal type selection

Open Topics



- Multi dimensional arrays - Delinearization
- Integer modulo arithmetics
- Optimal type selection



Polly alone is not yet improving performance of your code and may even apply incorrect transformations on the LLVM-IR.

Thanks to Qualcomm for sponsoring this talk.

Thank you.
Any Questions?

<http://wiki.llvm.org/Polly>