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## Tensor Evolution

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## Tensor Evolution

- Extension of LLVM Scalar Evolution (SCEV) for Tensors
- Analysis and Optimization Technique
- Tensors are
- multi-dimensional arrays
- fundamental to Machine Learning models


## Scalar Evolution (SCEV)

"Scalar Evolution is an LLVM analysis that is used to analyze, categorize and simplify expressions in loops. Many optimizations such as - generalized loop-strength-reduction, parallelization by induction variable (vectorization), and loop-invariant expression elimination - rely on SCEV analysis.
However, SCEV is also a complex topic."
-- some Large Language Model


## The LLVM Compiler Infrastructure

2018 European LLVM Developers Meeting
Scalar Evolution - Demystified
J. Absar

This is a tutorialtechnical-talk proposal for an illustrative and in-depth exposition of Scalar Evolution in LLVM. Scalar Evolution is an L analyse, categorize and simplify expressions in loops. Many optimizations such as - generalized loop-strength-reduction, parallelisation $t$ (vectorization), and loop-invariant expression elimination - rely on SCEV analysis.
However, SCEV is also a complex topic. This tutorial delves into how exactly LLVM performs the SCEV magic and how it can be used t analyse different optimisations.
This tutorial will cover the following topics:

1. What is SCEV? How does it help improve performance? SCEV in action (using simple clear examples), hain of Recurrences - which forms the mathematical basis of SCEV.
Simplifyyngrewriting rules in CR that SCEV uses to simplify expressions evolving out of induction variables. Terminology and SC AddRec) that is common currency that one should get familiar with when trying to understand and use SCEV in any context.
2. How to use SCEV analysis to write your own optimisation pass? Usage of SCEV by LSR (Loop Strength Reduce) and others 6. How to generate analysis info out of SCEV and how to interpret them

The last talk on SCEV was in LLVM-Dev 2009. This tutorial will be complementary to that and go further with examples, discussions an Ilvm since then. The author has previously given a talk on machine scheduler in llvm - https://www.youtube.com/watch?v=brpomKUyyE

## Scalar Evolution

- SCEV analysis and opt

```
int foo(int *a, int n, int k){
    for (int i=0; i < n; i++)
    a[i] = i*k;
}
```

\$ opt -analyze -scalar-evolution foo.II

1. Printing analysis 'Scalar Evolution Analysis' for function 'foo':
2. Classifying expressions for: @foo
3. 
4. $\%$ mul $=$ mul nsw $i 32 \%$ i, \%k
5. --> $\{0,+, \% k\}<\%$ for body> Exits: $((-1+\% n) * \% k)$
6. 

## Tensor Evolution - Motivating Example 1

```
# PyTorch code.
# a and x are tensors
def forward(self, a, x):
    for _ in range(15):
        x = a + x
    return x
```

- Tensor Evolution Optimization

```
# PyTorch code.
# a and x are tensors
def forward(self, a, x):
    return 15*a+x
```


## Mathematical Formulation

- Basic Recurrence (Tensor Evolution)
- a constant or loop-invariant tensor $\mathrm{T}_{\mathrm{c}}$
- a function $\tau_{1}$ over natural number $N$ that produces tensor of same shape as $T_{c}$
- an element-wise operator + associative and commutative
- $\tau$ defined as function $\tau$ (i) over $N$

$$
\begin{aligned}
& \tau=\left\{T_{c},+, \tau_{1}\right\} \\
& \left\{T_{c^{\prime}}+, \tau_{1}\right\}(i)=T_{c}+\tau_{1}(0)+\tau_{1}(1) \ldots+\tau_{1}(i-1)
\end{aligned}
$$

## Mathematical Formulation

- Chain of Recurrences (Tensor Evolution)
- loop-invariant tensors $\mathrm{Tc}_{0}, \mathrm{Tc}_{1}, \mathrm{Tc}_{2}, \ldots, \mathrm{Tc}_{\mathrm{i}-1}$;
- function $\tau_{k}$ defined over $N$,
- operators $\bigodot_{1}, \bigodot_{2}, \ldots, \bigodot_{k}$,
- chain of evolution of tensor value represented by tuple

$$
\begin{array}{ll}
\tau=\left\{T c_{0}, \bigodot_{1}, T c_{1}, \bigodot_{2}, \ldots, \bigodot_{k}, \tau_{k}\right\} & \text { eq. } 1 \\
\tau(i)=\left\{T c_{0}, \bigodot_{1},\left\{T c_{1}, \bigodot_{2}, \ldots, \bigodot_{k}, \tau_{k}\right\}\right\}(i) & \text { eq. } 2
\end{array}
$$

- Note: Operators could be same or different (+,-, *, tanh).
- Recurrences
- Algebraic properties
- Computationally reducible at any iteration point


## Tensor Evolution

- Lemmas - Rewrite Rules
- Used for building TEV 'available' expressions and simplifications

| operator | TEV expression | rewrite rule |
| :--- | :--- | :--- |
| slice | slice $(\{A,+, \tau\})$ | $\{\operatorname{slice}(A),+$, slice $(\tau)\}$ |
|  | slice $(\{A, *, \tau\})$ | $\{\operatorname{slice}(A), *$, slice $(\tau)\}$ |
| reshape | $\operatorname{reshape}(\{A, \odot, \tau\})$ | $\{\operatorname{reshape}(A), \odot, \operatorname{reshape}(\tau)\}$ |
| concat | $\operatorname{concat}\left(\left\{A, \odot, \tau_{1}\right\},\left\{B, \odot, \tau_{2}\right\}\right)$ | $\left\{\operatorname{concat}(A, B), \odot, \operatorname{concat}\left(\tau_{1}, \tau_{2}\right)\right\}$ |
| add K | $K+\{A,+, \tau\}$ | $\{K+A,+, \tau\}$ |
| add TEVs | $\left\{A,+, \tau_{1}\right\}+\left\{B,+, \tau_{2}\right\}$ | $\left\{A+B,+, \tau_{1}+\tau_{2}\right\}$ |
| mul | $K *\{A,+, \tau\}$ | $\{K * A,+, K * \tau\}$ |
| inject TEV | $\{A,+,\{B,+, \tau\}\}$ | $\{A,+B,+, \tau\}$ |

## Tensor Evolution - Basic Recurrence



## Tensor Evolution

- Lemma: Add a constant (LIV) tensor

loop-exit


## Tensor Evolution

- Lemma: Add two TEVs

> loop-header

loop-exit

## Tensor Evolution

- Lemma: TEV inject into TEV



## Tensor Evolution

- Lemma: Slice

loop-exit


## Tensor Evolution

- Lemma: Reshape
loop-header

loop-exit


## Tensor Evolution

- Lemma: Concat

loop-exit


## Tensor Evolution

- Lemmas - Rewrite Rules
- Used for building TEV expressions and simplifications

| operator | TEV expression | rewrite rule |
| :--- | :--- | :--- |
| slice | slice $(\{A,+, \tau\})$ | $\{\operatorname{slice}(A),+$, slice $(\tau)\}$ |
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| add K | $K+\{A,+, \tau\}$ | $\{K+A,+, \tau\}$ |
| add TEVs | $\left\{A,+, \tau_{1}\right\}+\left\{B,+, \tau_{2}\right\}$ | $\left\{A+B,+, \tau_{1}+\tau_{2}\right\}$ |
| mul | $K *\{A,+, \tau\}$ | $\{K * A,+, K * \tau\}$ |
| inject TEV | $\{A,+,\{B,+, \tau\}\}$ | $\{A,+, B,+, \tau\}$ |

TEV Pass - Analysis


## TEV Pass - Opt



```
# PyTorch code.
def forward(self, a, x, y):
    for _ in range(15):
        x = x + a
        z = x[1,:]
        y = y + z
    return y
```


## Evaluation of $\mathrm{Y}_{\underline{k}}$

$Y_{k}=\left\{Y_{\theta},+, S\left(\left\{X_{0},+, A\right\}\right)\right\}_{k}$
$\rightarrow Y_{k}=\left\{Y_{\theta},+, S\left(\left\{X_{\theta},+, A\right\}\right)\right\}_{k}$
$\rightarrow Y_{k}=\left\{Y_{\theta},+,\left\{S\left(X_{\theta}\right),+, S(A)\right\}\right\}_{k}$
$\rightarrow Y_{k}=\left\{Y_{\theta},+, S\left(X_{\theta}\right),+, S(A)\right\}_{k}$
$\rightarrow Y_{k}=Y_{\theta}+k * S\left(X_{\theta}\right)+k^{*}(k+1) / 2 * S(A)$

## TEV Pass - Opt

> Evaluation of $Y_{\underline{k}}$
> $Y_{k}=\left\{Y_{\theta},+, S\left(\left\{X_{\theta},+, A\right\}\right)\right\}_{k}$
> $\rightarrow Y_{k}=\left\{Y_{\theta},+, S\left(\left\{X_{\theta},+, A\right\}\right)\right\}_{k}$
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> $\rightarrow Y_{k}=Y_{\theta}+k^{*} S\left(X_{\theta}\right)+k^{*}(k+1) / 2^{*} S(A)$

```
# PyTorch code.
def forward(self, a, x, y):
    for _ in range(15):
        x = x + a
        z = x[1,:]
        y = y + z
    return y
```

```
# PyTorch code.
def forward(self, a, x, y):
    return y + 15*x[1,:] + 15*(15+1)/2*a[1,:]
```


## Conclusion

- TEV is extension of SCEV to Tensors
- Construction of TEV expressions and rewrite-lemmas
- Complex optimizations on top of TEV (much like SCEV LSR etc)
- Prototyped in internal-compiler
- Potential opt for MLIR lower CFG dialects
- Looking forward to collaboration and discussions


## Thank you

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