

Precise and Efficient Garbage Collection in VMKit with MMTk

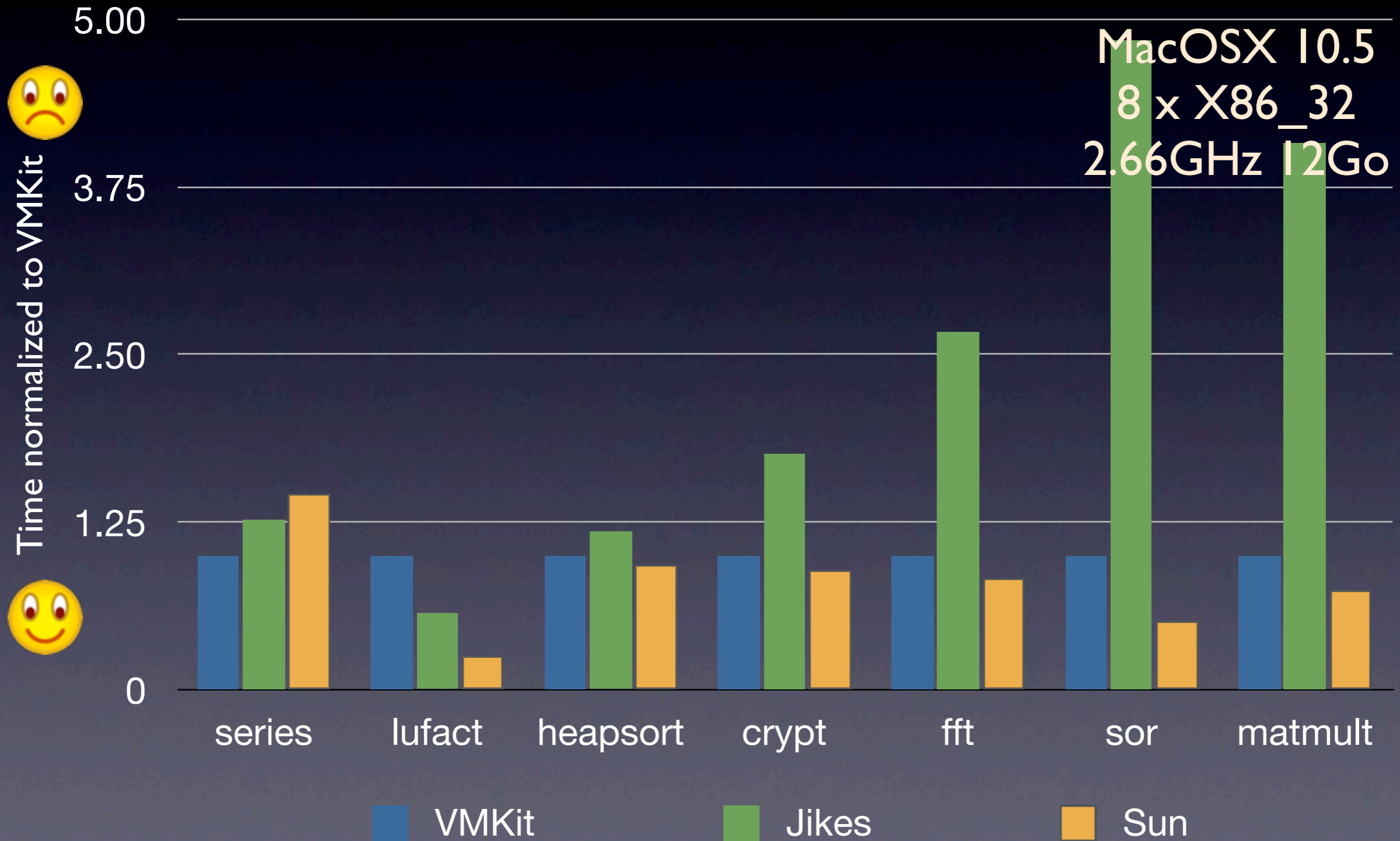
LLVM Developer's Meeting
Nicolas Geoffray

nicolas.geoffray@lip6.fr

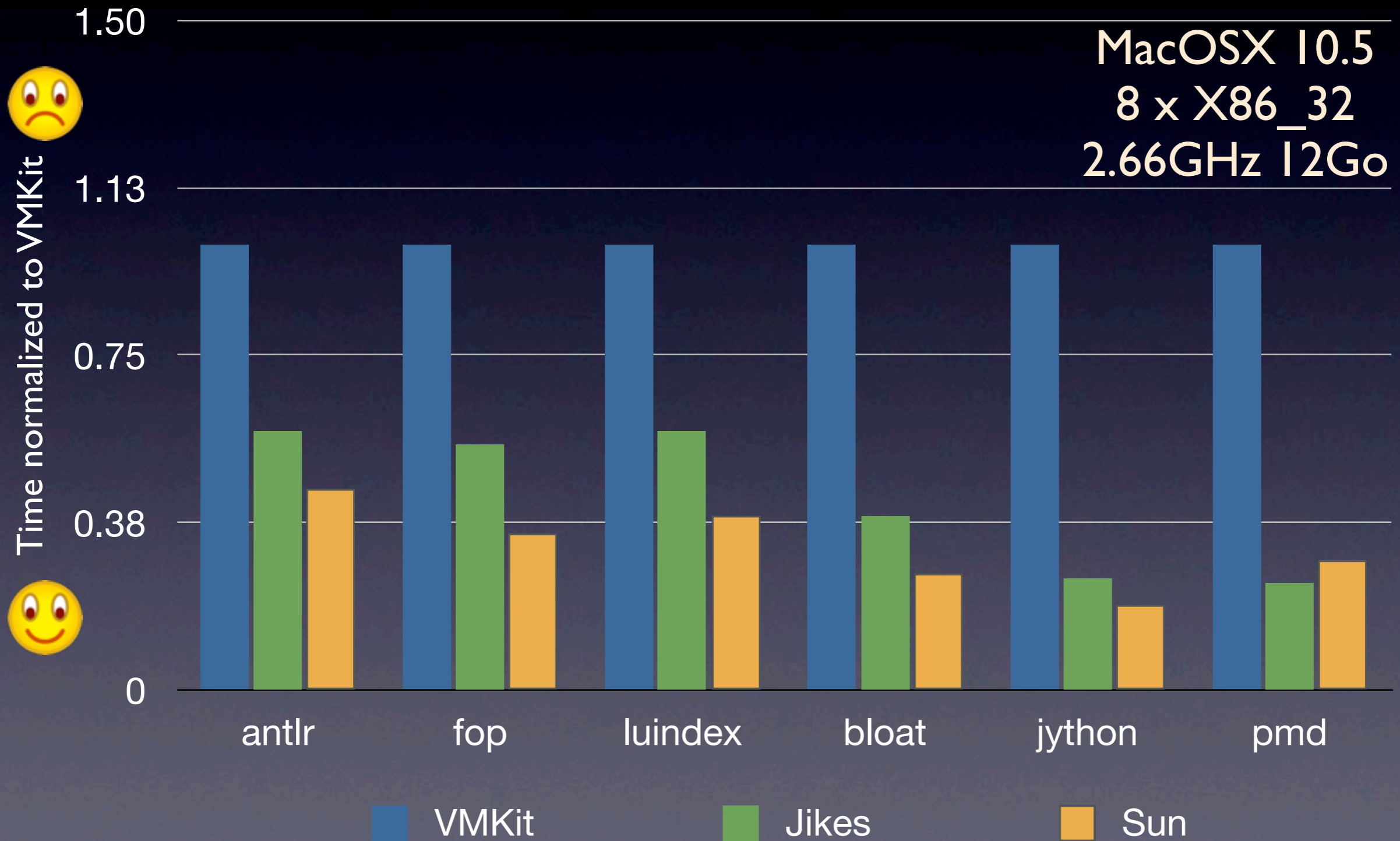
Background

- VMKit: Java and .Net on top of LLVM
 - Uses LLVM's JIT for executing code
 - Uses Boehm for GC
- Performance bottlenecks
 - No dynamic optimization
 - Conservative GC

CPU-intensive Benchmarks (JGF)

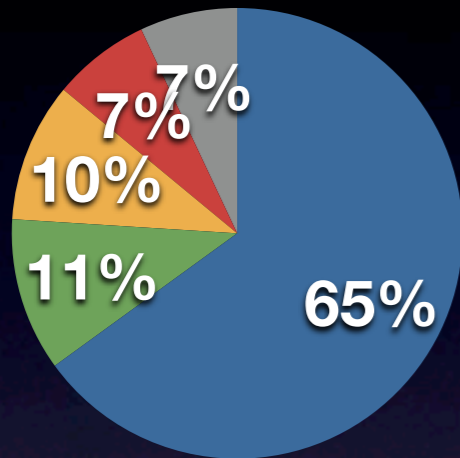


VM-intensive Benchmarks (Dacapo)

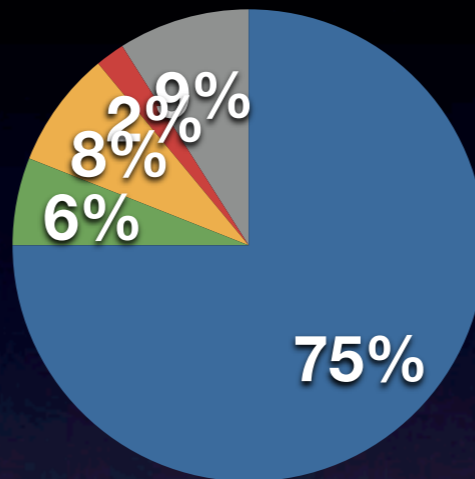


Execution Overheads

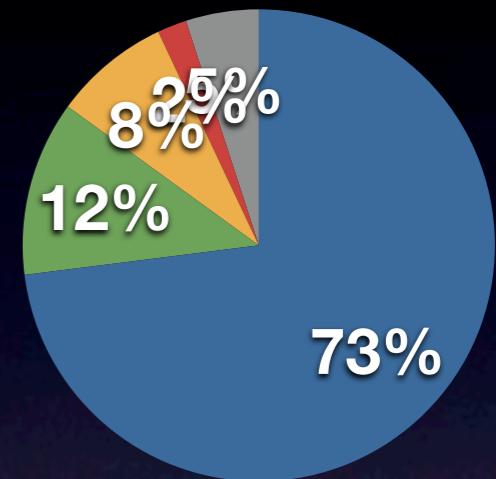
antlr



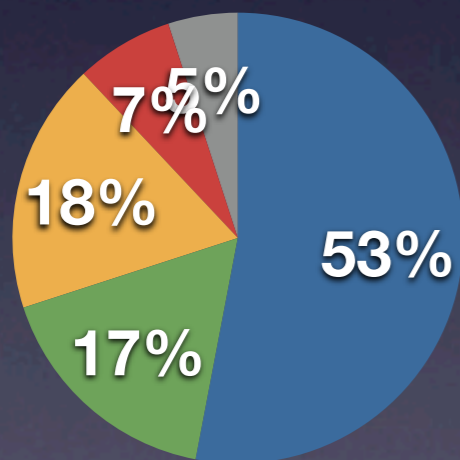
fop



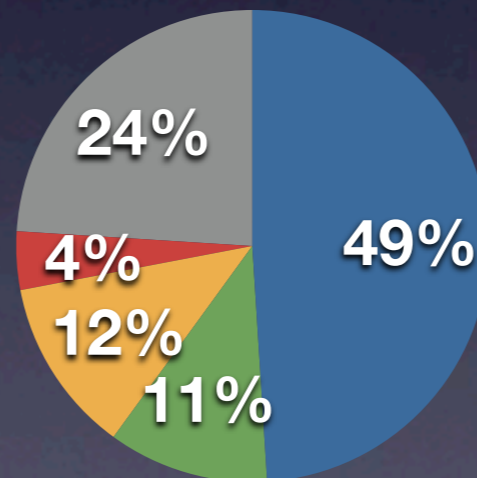
luindex



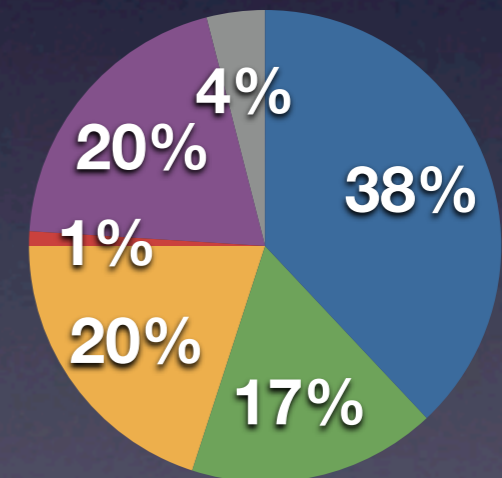
bloat



jython



pmd



Application

Allocations

Collections

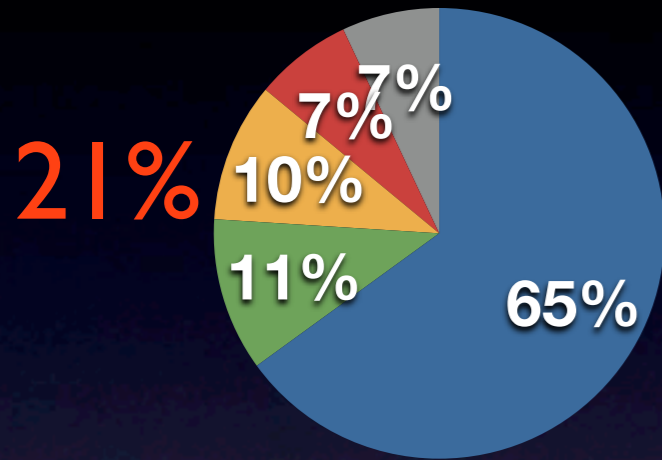
System.arraycopy

Interface calls

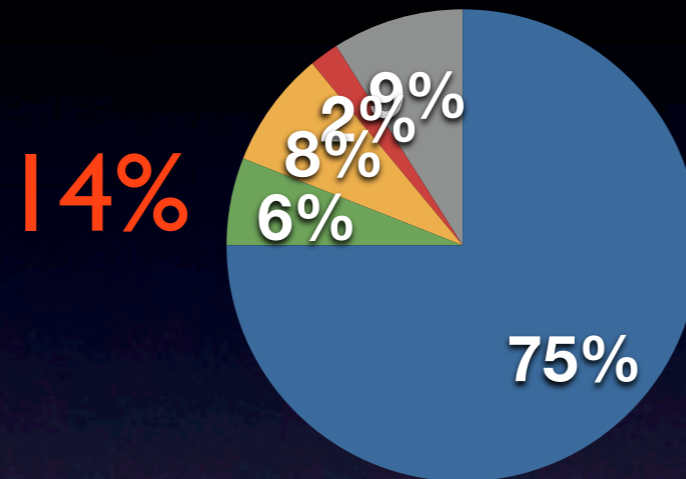
Others

Execution Overheads

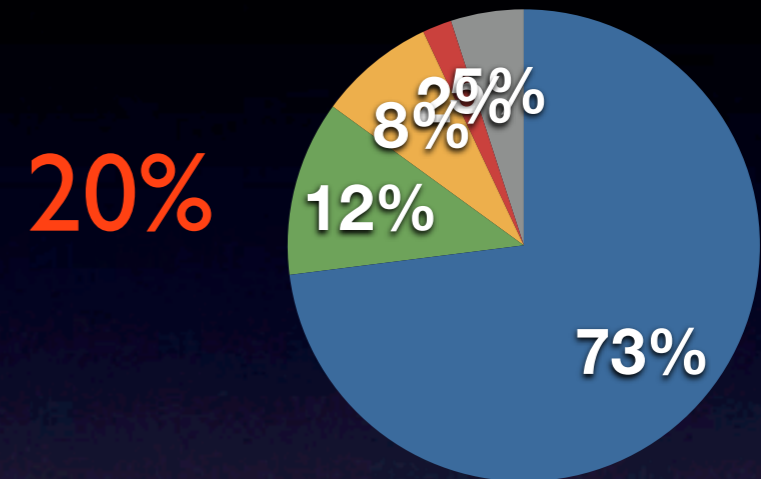
antlr



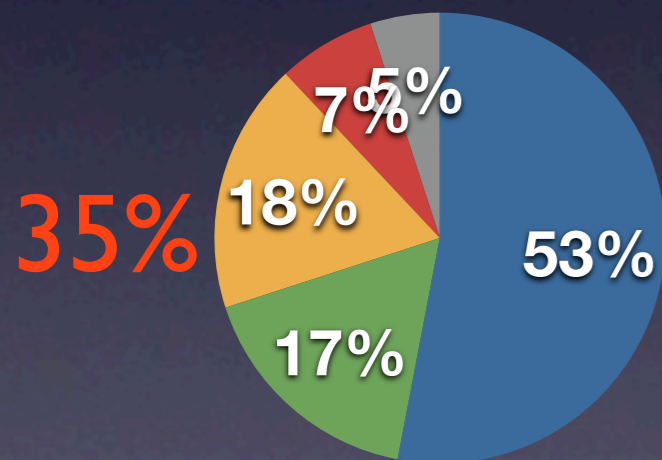
fop



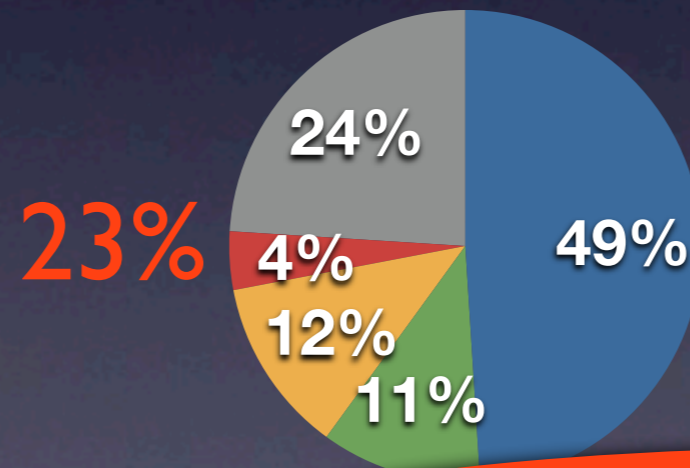
luindex



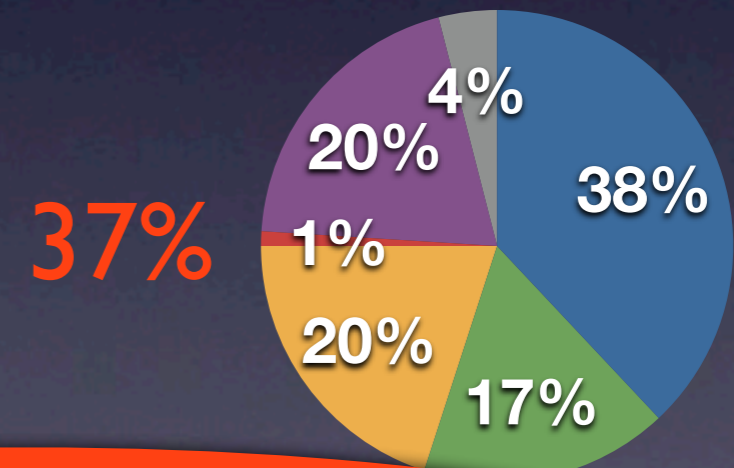
bloat



jython



pmd



Application

System.arraycopy

Allocations

Interface calls

Collections

Others

Goal: replace Boehm with MMTk

- MMTk is JikesRVM's GC
 - Framework for writing GCs
 - Multiple GC Implementations (Copying, Mark and trace, Immix)
- Copying collectors require precise stack scanning
 - Locate pointers on the stack

But... it's in Java?

- Yes, but nothing to be afraid of:
 - Use of **Magic** tricks
 - No use of runtime features (exceptions, inheritance)
 - No use of standard library
- Use VMKit's AOT compiler
 - Transform MMTk into a .bc file

Outline

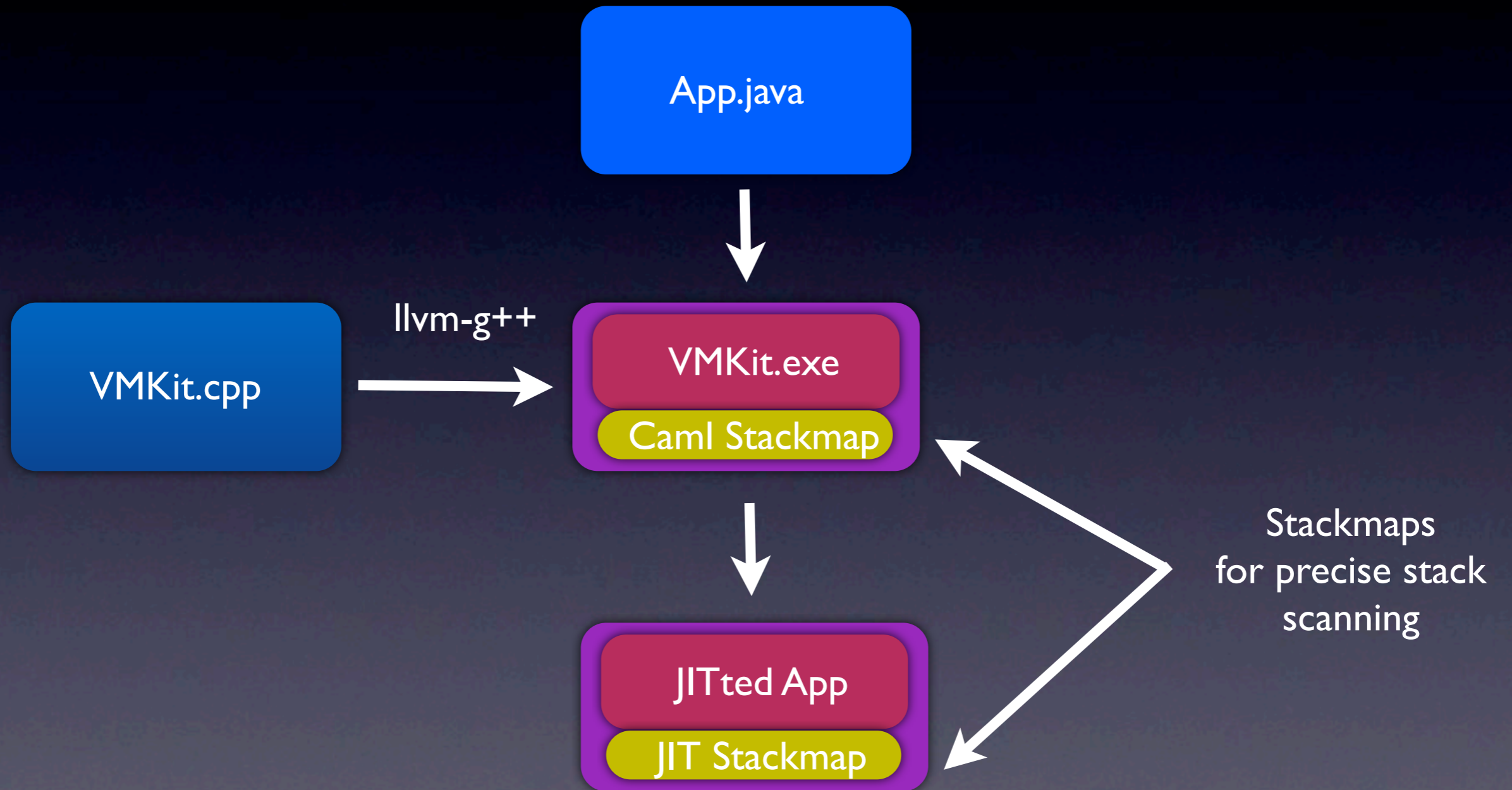
- Introduction
- Precise garbage collection
- Compiling MMTk with VMJC
- Putting it all together
- What's left

- Introduction
- **Precise garbage collection**
- Compiling MMTk with VMJC
- Putting it all together
- What's left

Precise Garbage Collection

- Write code that locates pointers in the stack
 - `llvm.gcroot` in JIT-generated code
 - `llvm.gcroot` in VMKit's runtime written in C++
- Use LLVM's GC framework to generate stack maps
 - Caml stack maps for `llvm-g++` generated code
 - JIT stack maps for JIT-generated code

Precise Garbage Collection



Stack Scanning

- Problem: interweaving of different kinds of functions
 - Application's managed (Java or C#) functions: **trusted**
 - VMKit's C++ functions: **trusted**
 - Application's JNI functions: **untrusted**
- Solution: create a side-stack for frame addresses
 - Updated upon entry of a kind of method
 - VMKit knows the kind of each frame on the thread stack

Type of methods

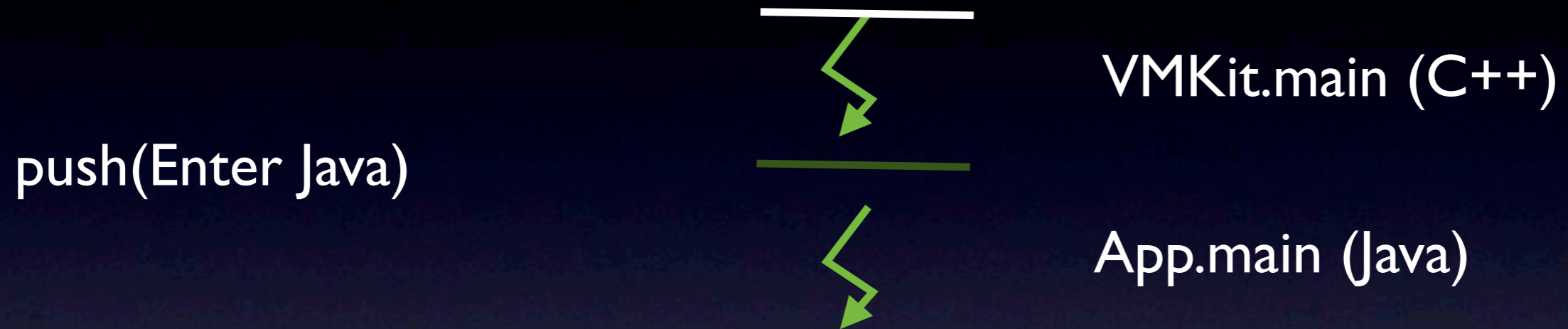
- **Trusted**
 - Has a stack map, so can manipulate objects (llvm.gcroot)
 - Saves frame pointer (llvm::NoFramePointerElim)
- **Untrusted**
 - Has no stack map, so should not manipulate objects
 - May not save the frame pointer

Stack Scanning Example (I)



VMKit.main (C++)

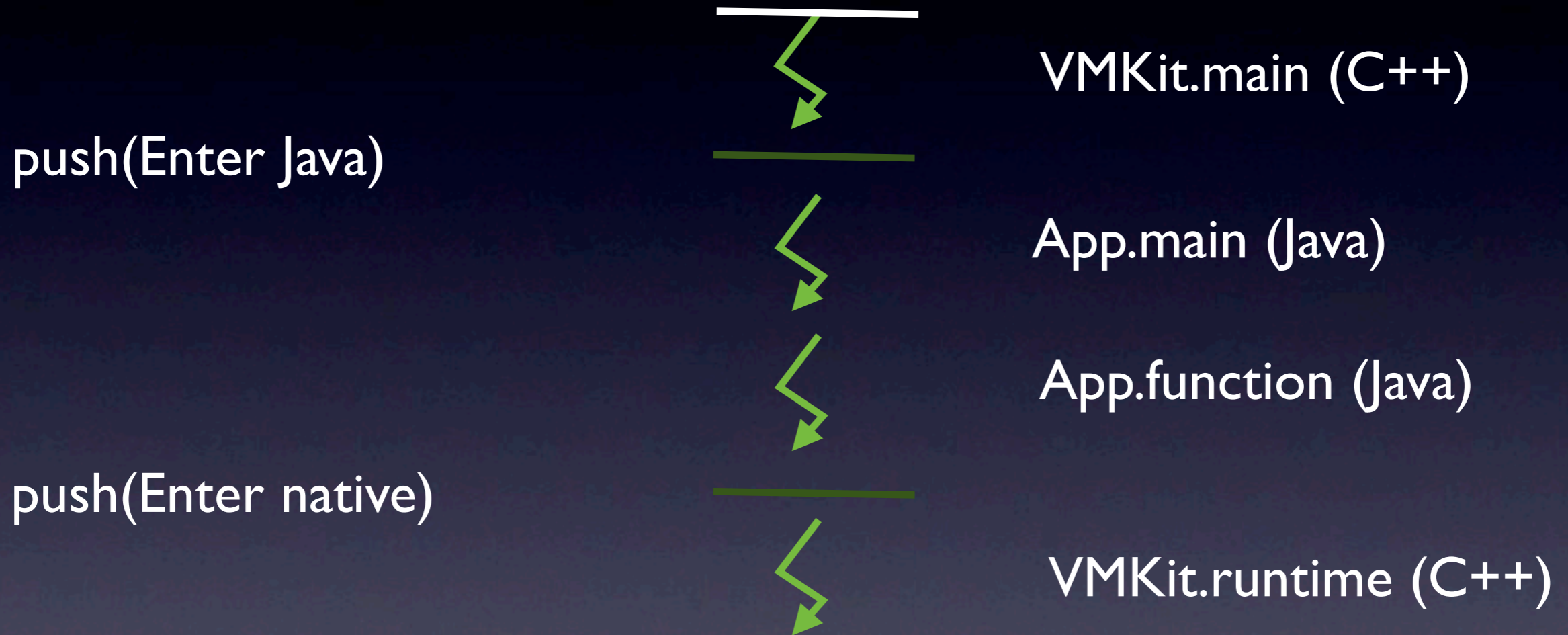
Stack Scanning Example (I)



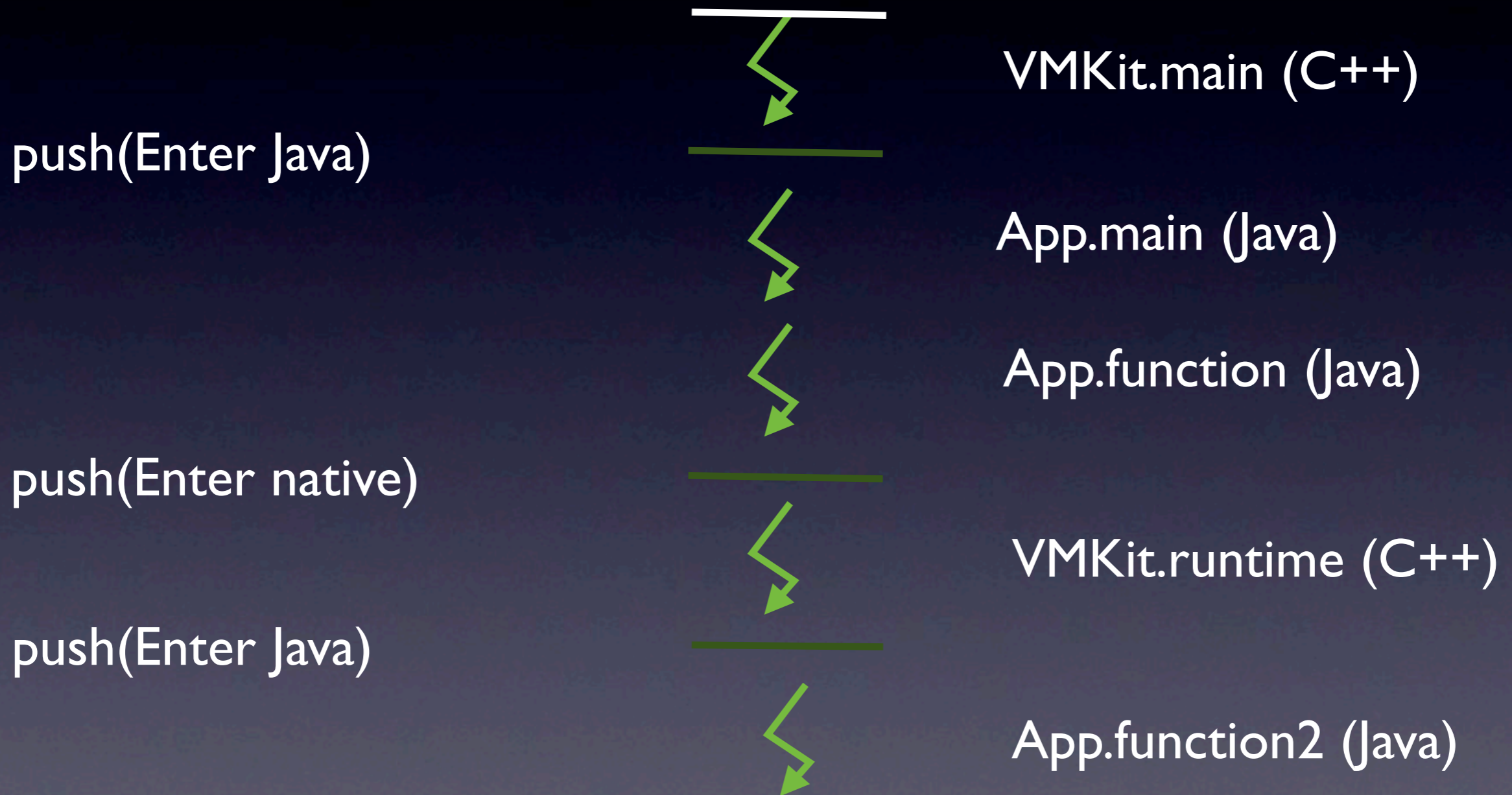
Stack Scanning Example (I)



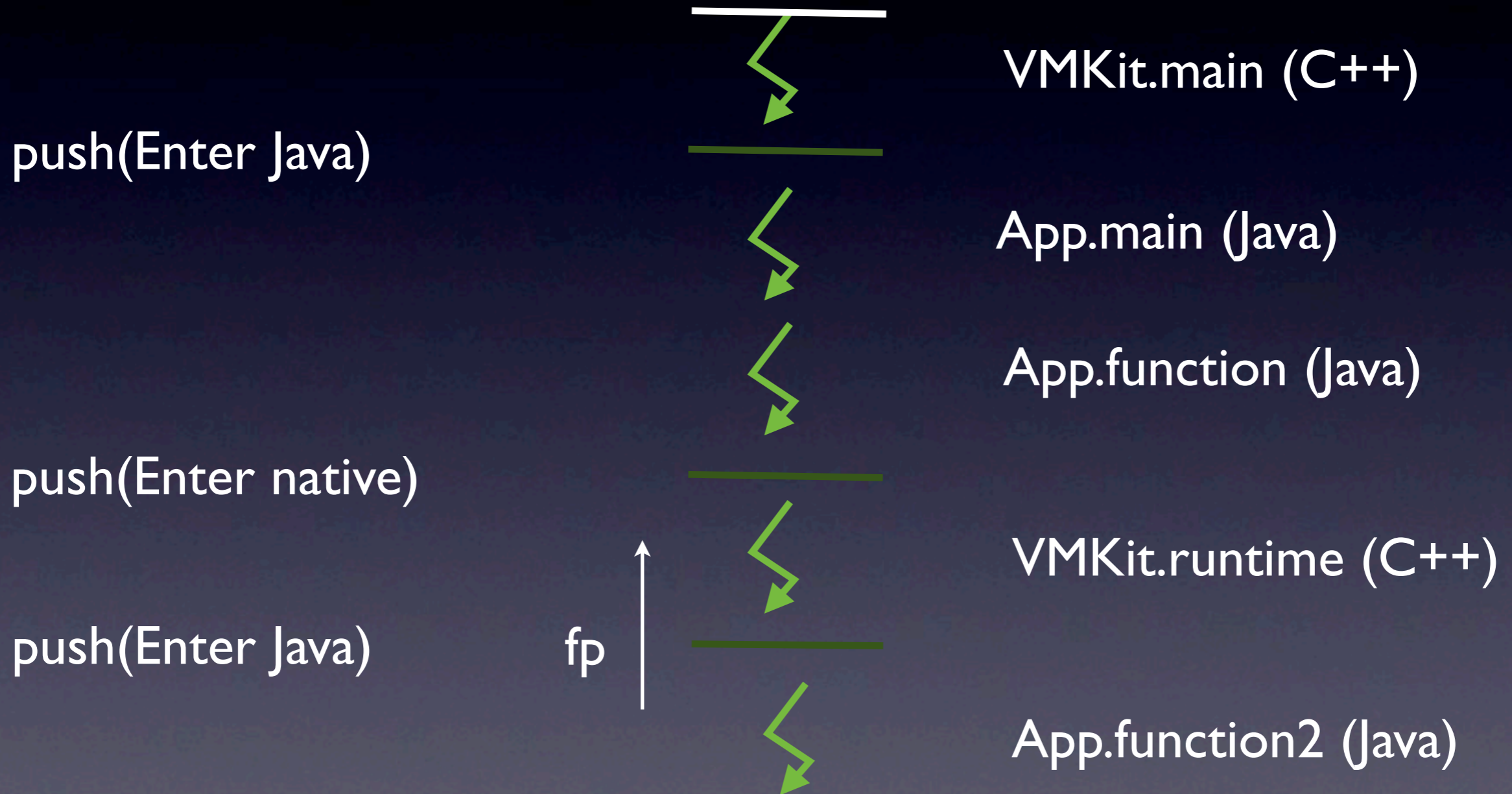
Stack Scanning Example (I)



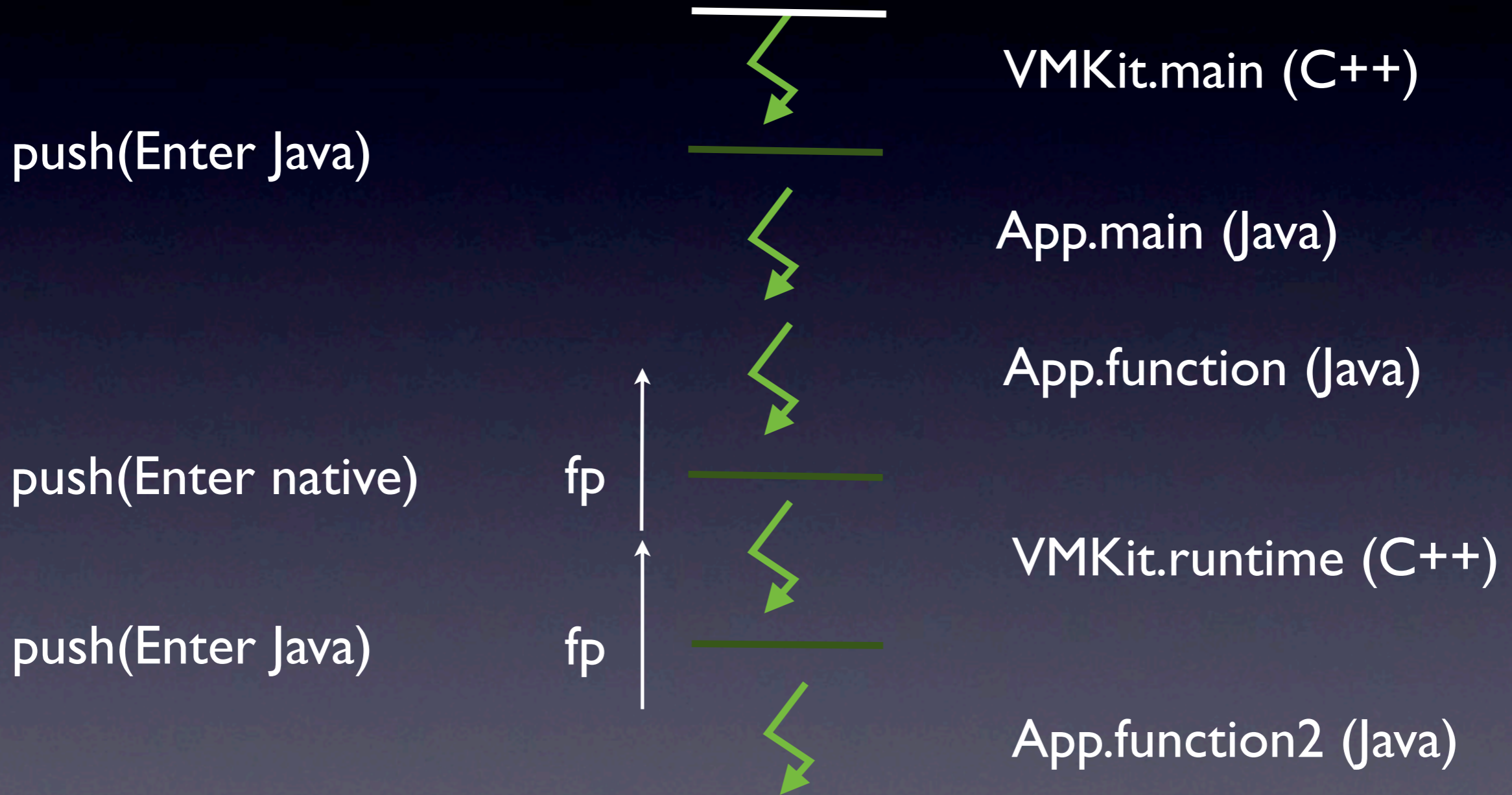
Stack Scanning Example (I)



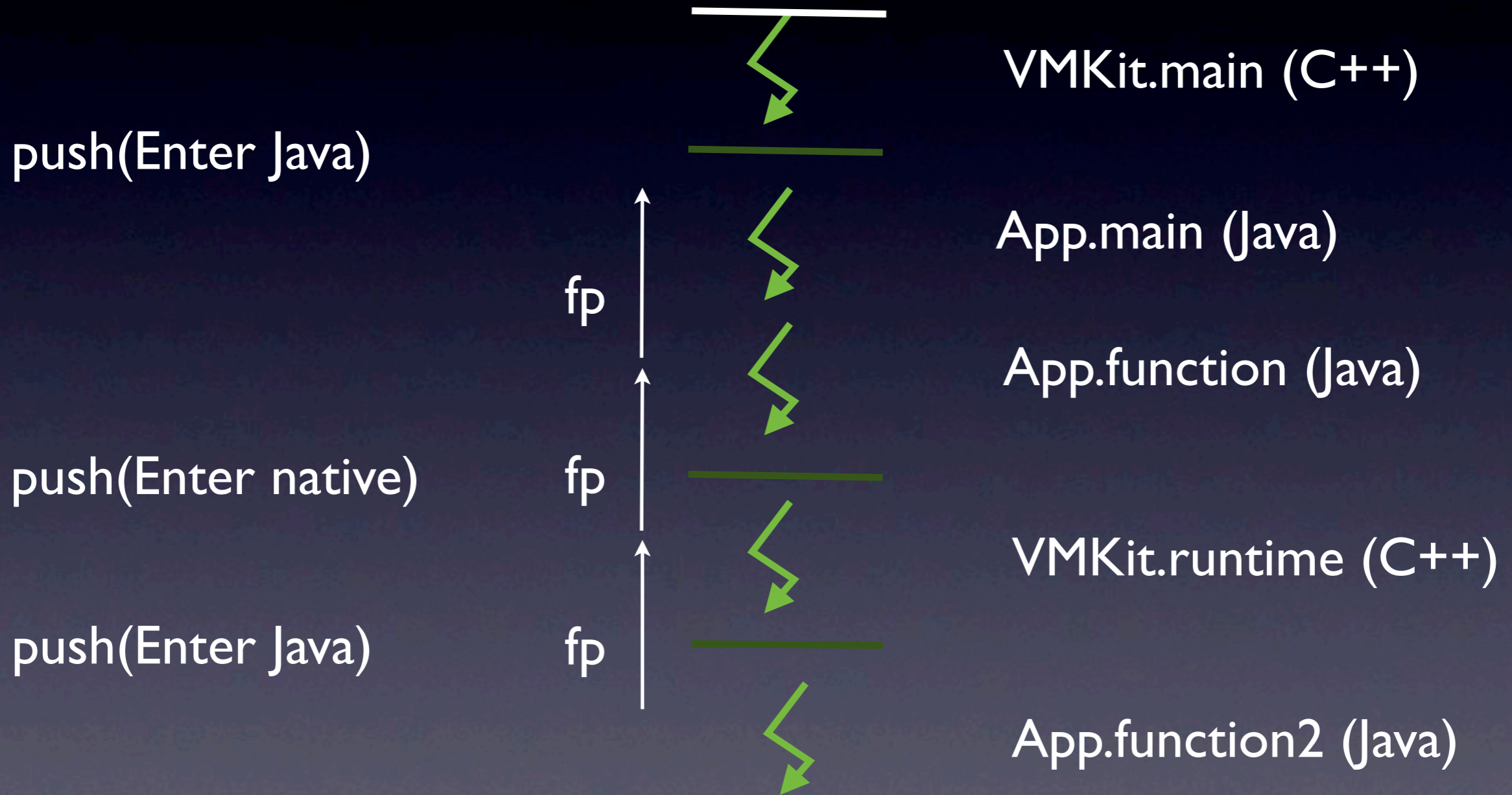
Stack Scanning Example (I)



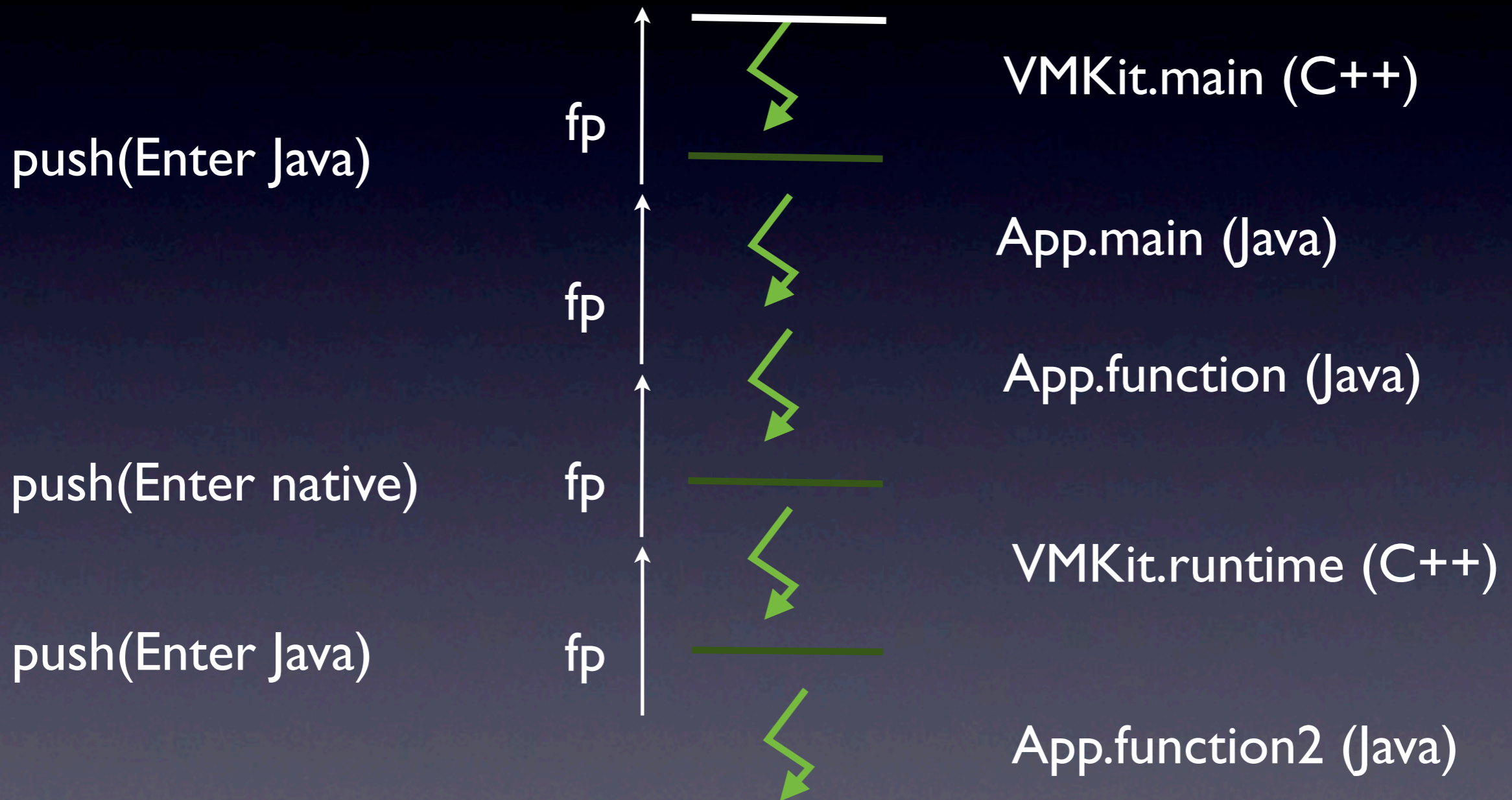
Stack Scanning Example (I)



Stack Scanning Example (I)



Stack Scanning Example (I)

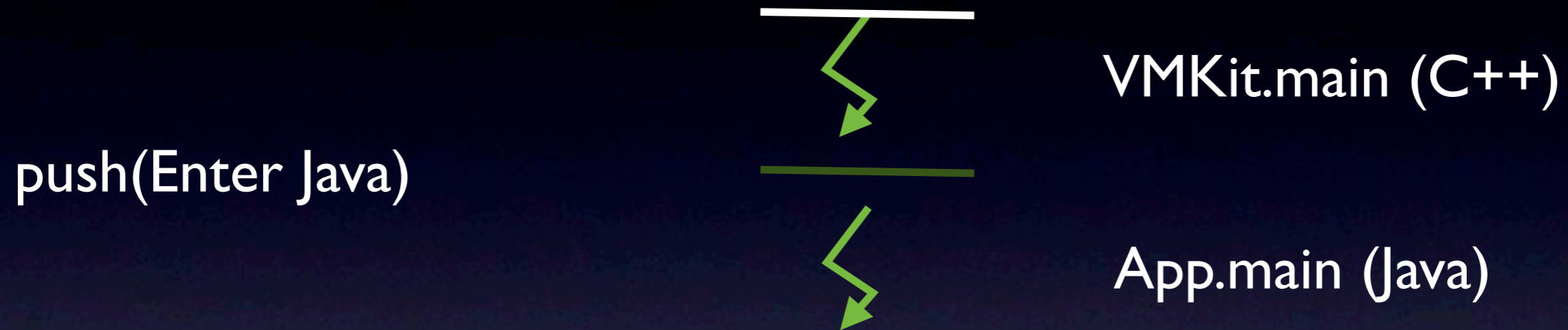


Stack Scanning Example (2)

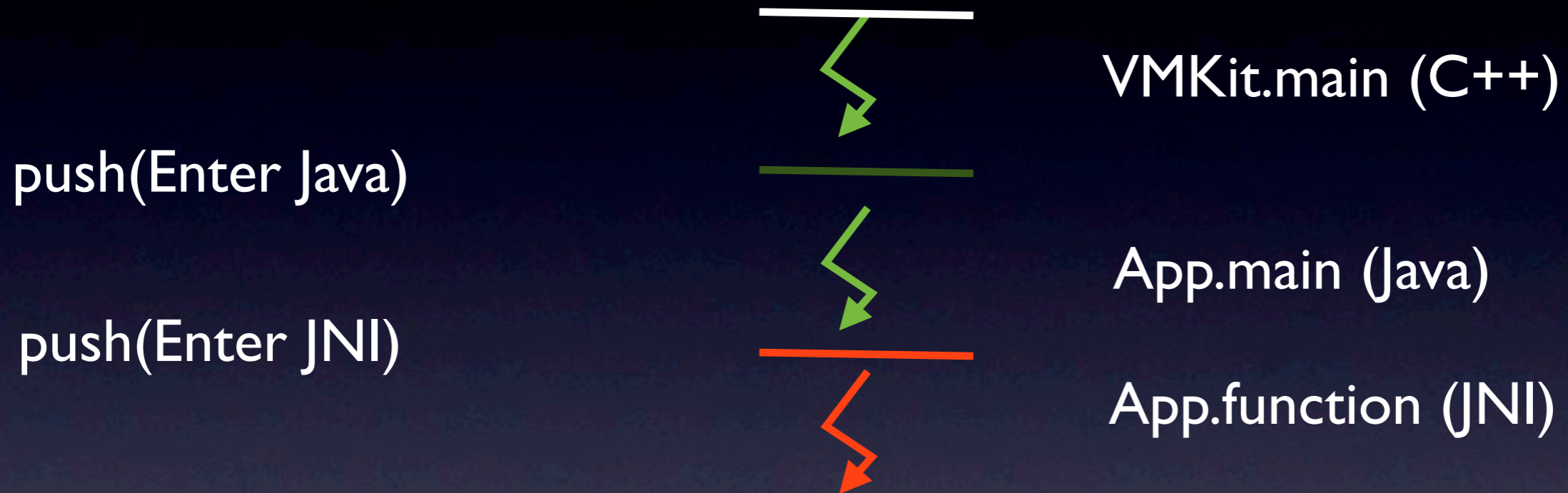


VMKit.main (C++)

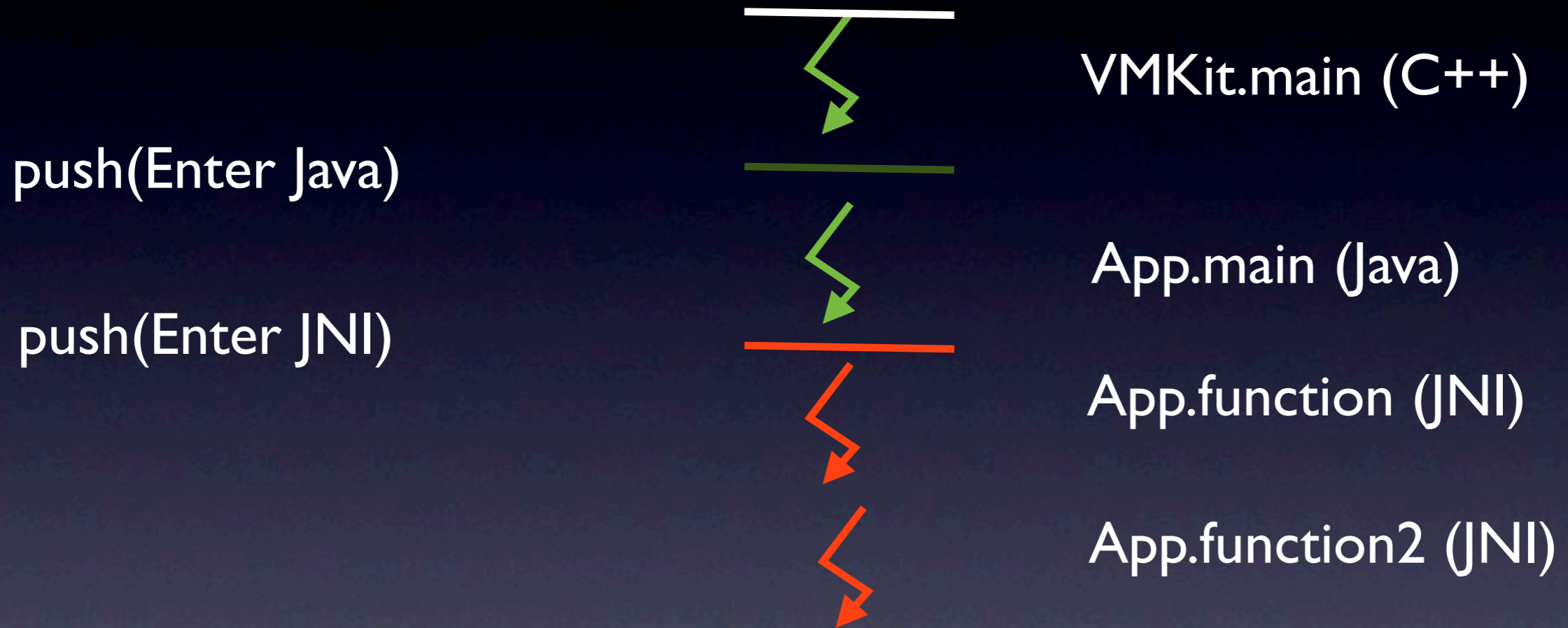
Stack Scanning Example (2)



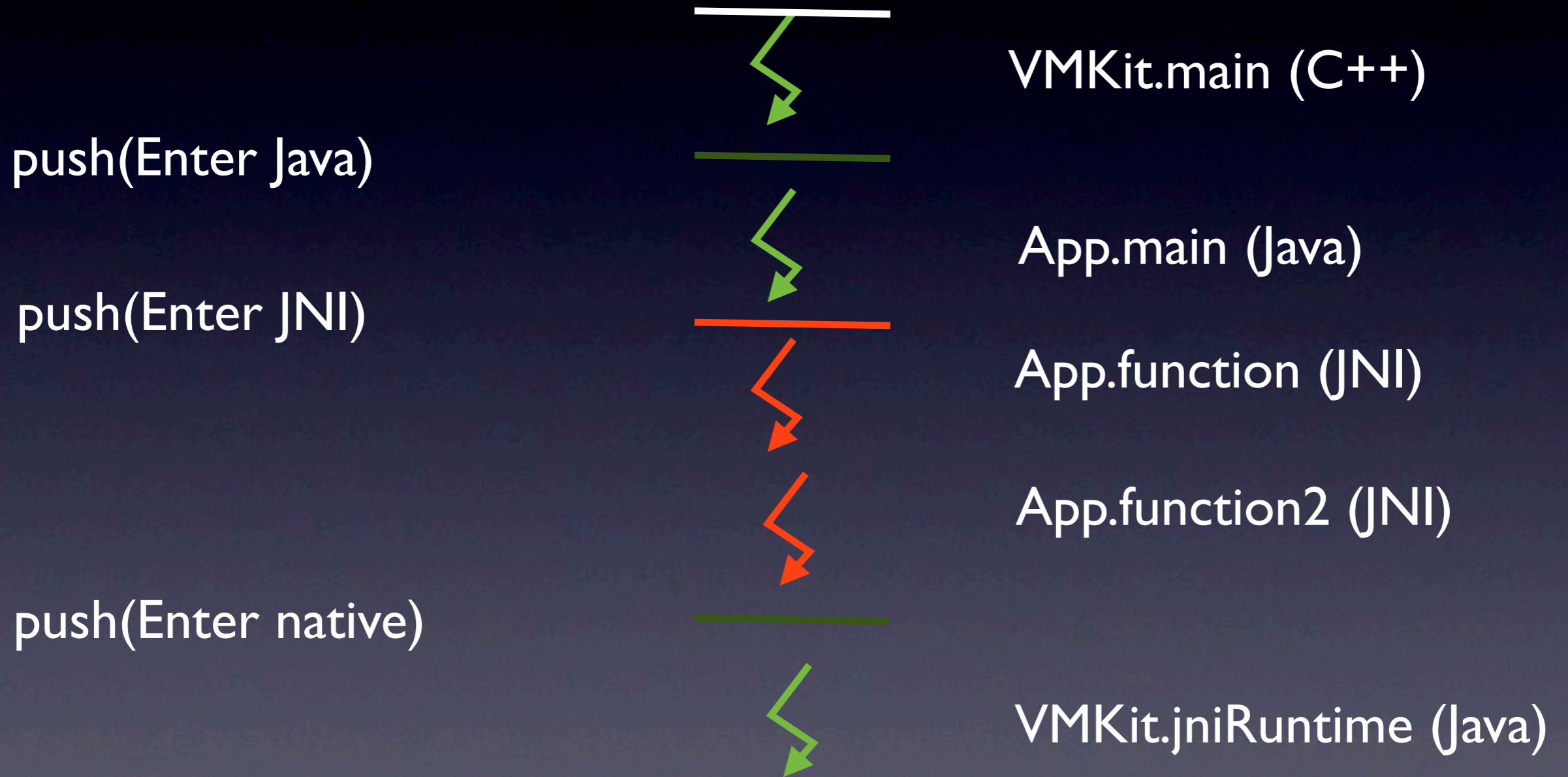
Stack Scanning Example (2)



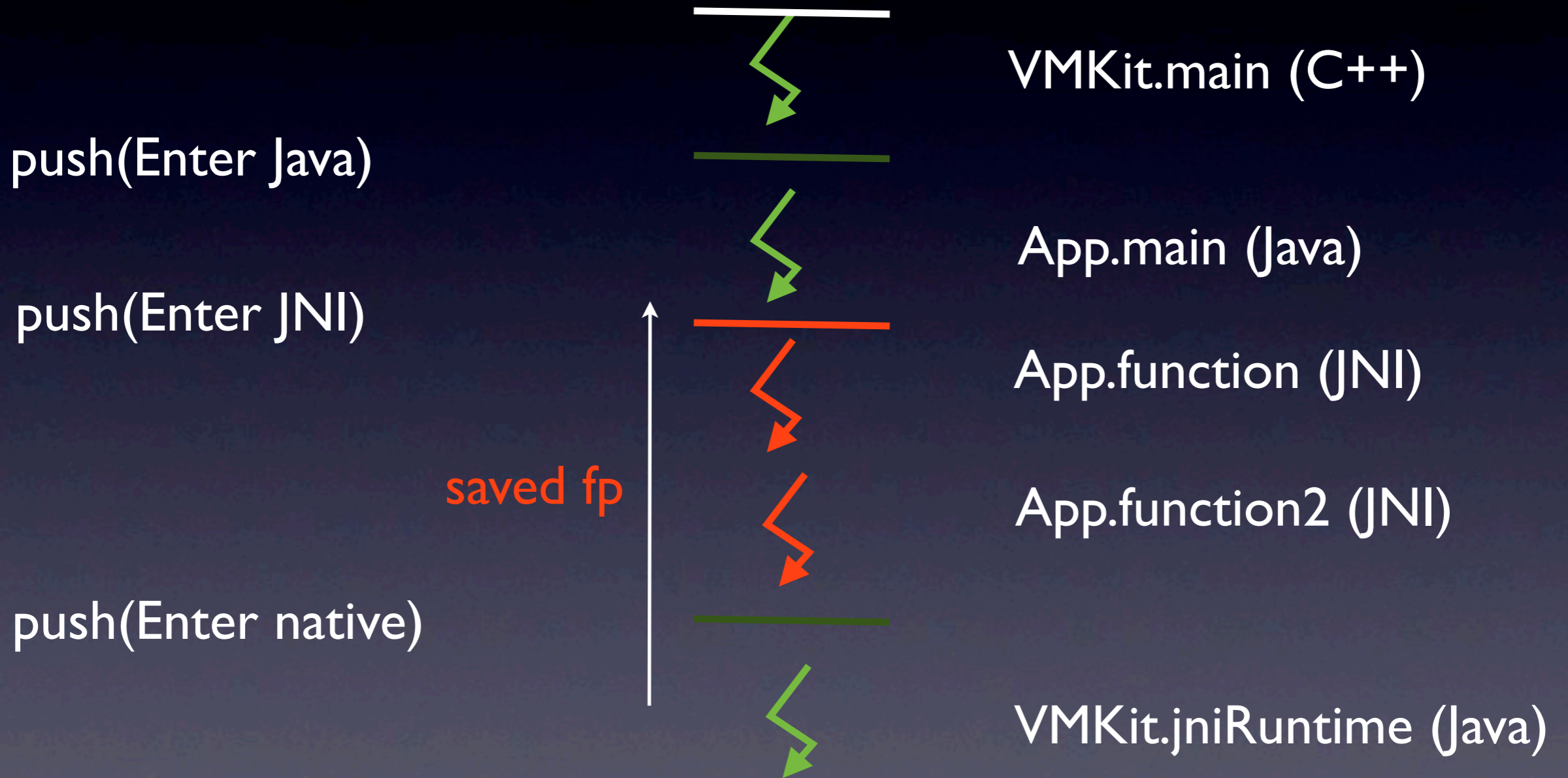
Stack Scanning Example (2)



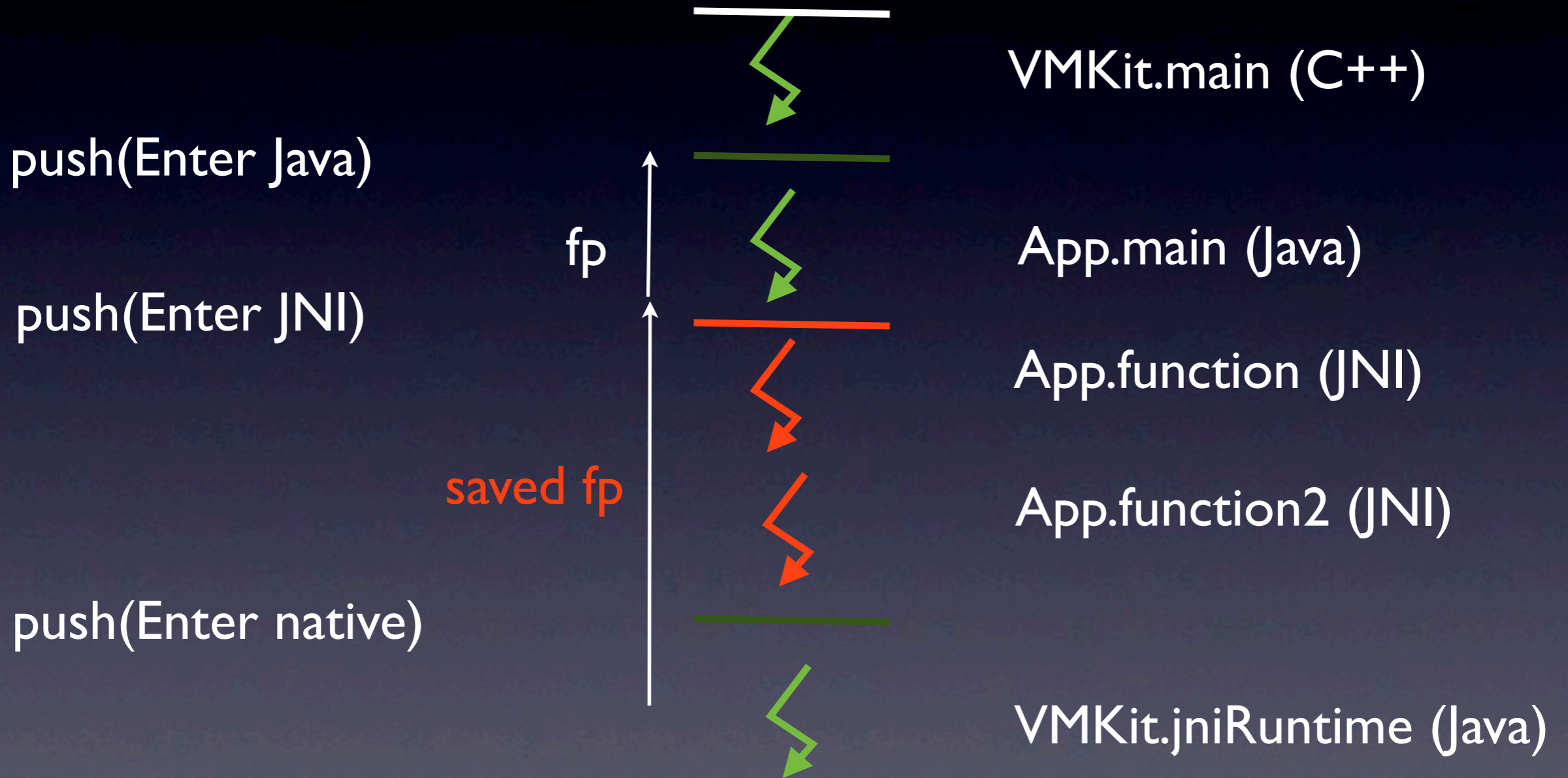
Stack Scanning Example (2)



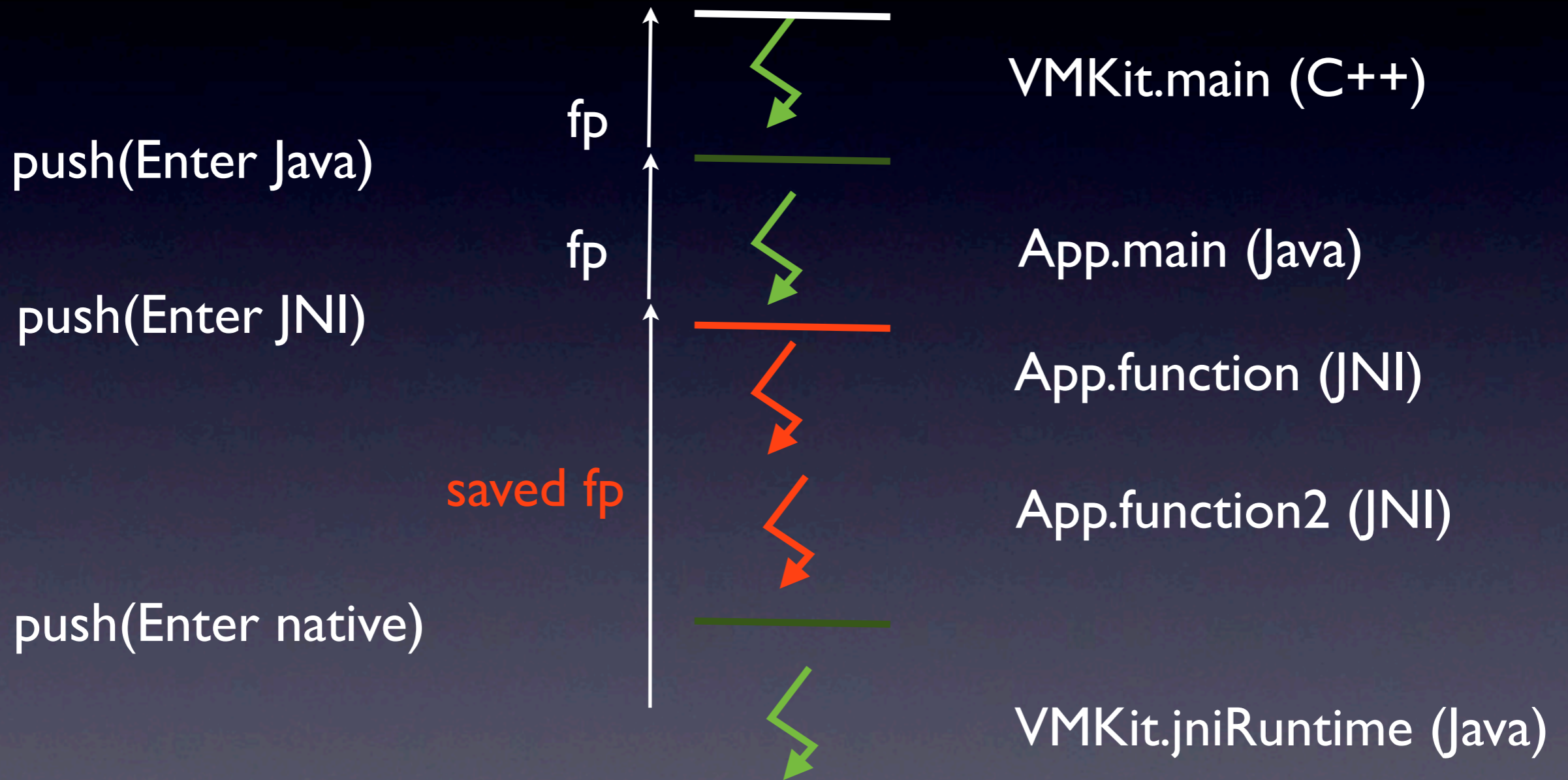
Stack Scanning Example (2)



Stack Scanning Example (2)



Stack Scanning Example (2)



Running the GC

A precise GC scans the stacks at *safe points*: point during execution where the GC can know the type of each value on the stack

Single-threaded Application

- GC always triggered at *safe points*
 - `gmalloc` instructions
 - `Collector::collect()`

Multi-threaded Application

- When entering a GC, must wait for all threads to join
 - Don't use signals! or no safe point
 - Use a thread-local variable to poll on method entry and backward branches
 - Scan stacks of threads blocked in JNI or system calls

Application changes for GC

```
public static void runLoop(int a) {  
    while (a-- > 0) System.out.println("Hello World");  
}
```

Application changes for GC

```
public static void runLoop(int a) {  
    if (getThreadID().doGC) GC()  
    while (a-->0) {  
        System.out.println("Hello World");  
        if (getThreadID().doGC) GC()  
    }  
}
```

- Introduction
- Precise garbage collection
- **Compiling MMTk with VMJC**
- Putting it all together
- What's left

What is VMJC?

- An Ahead of Time compiler (AOT)
 - Generates .bc files from .class files
- Use of llvm tools to generate platform-dependant files
 - shared library: `llc -relocation-model=pic + gcc`
 - executable: `llc + ld vmkit + gcc`

Goal: compile MMTk with VMJC

- Generate a .bc file that can be linked with VMKit
 - Interface MMTK → VMKit (e.g. threads synchronization, stack scanning)
 - Interface VMKit → MMTk (e.g. gcmalloc)

Why MMTk does not need a Java runtime?

- No use of runtime features
 - synchronizations, exceptions, inheritance
- No use of standard library
 - HashMap, LinkedList, ArrayList

How MMTk is manipulating pointers?

- Definition of **Magic** classes and methods
 - Address, Word, Offset
 - Word Address.loadWord(Offset)
- **Magic** classes and methods translated by the compiler [VEE'09]
 - Similar mechanism than Inline ASM for C

Example (Frampton [VEE'09])

Inline ASM in C

```
void prefetchObjects(
    OOP *buffer,
    int size) {
    for(int i=0;i < size;i++){
        OOP o = buffer[i];
        asm volatile(
            "prefetchnta (%0)" ::
            "r" (o));
    }
}
```

Magic in Java

```
@NoBoundsCheck
void prefetchObjects(
    ObjectReference[] buffer) {
    for(int i=0;i<buffer.length;i++) {
        ObjectReference current
            = buffer[i];
        current.prefetch();
    }
}
```

- Introduction
- Precise garbage collection
- Compiling MMTk with VMJC
- Putting it all together
- What's left

Option 1: Object File

- Create a .o file of MMTk
 - `gcc mmtk.o vmkit.o -o vmkit`
- But...
 - No inlining in application code

Option 2: LLVM Bitcode File

- Create a .bc file of MMTk
 - `vmkit (-load mmtk.bc) -java HelloWorld`
- Late binding of allocations in VMKit code
 - `gcmalloc` in C++ are linked at runtime
- Inlining in Java code
 - `new` in applications are inlined with MMTk's `malloc`

Option 3: Everything is Bitcode

- Create a .bc file of MMTk
- Create a .bc file of VMKit
- Link, optimize and run

- Introduction
- Precise garbage collection
- Compiling MMTk with VMJC
- Putting it all together
- **What's left**

What's left

- Implementing the MMTK → VMKit interface
 - Interactions between the GC and the VM
- Finish implementation with read/write barriers
 - In VMKit code, in managed code
- Run benchmarks!
 - Benchmark with different GCs from MMTk

<http://vmkit.llvm.org>