
CIL : Common MLIR abstraction for C / C++/ Fortran

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Compiler Tree Technologies

CIL: Introduction and Status

- Middle level IR written as MLIR dialect.
- Comes in between AST and LLVM IR.
- Current focus is C / C++ / Fortran.
- Target / ABI independent.
- C/C++ input is parsed using clang AST plugin and converted to CIL.
- Fortran is parsed by FC compiler and converted to CIL.
- Working LTO Framework.
- Working Loop Nest Optimizer and Data Layout optimizations.
- SPEC CPU 2017 : Passes 2 C, 2 Fortran and 1 C++ benchmarks.

Why CIL?

- Flang is planning to use FIR, a MLIR dialect, for middle end optimizations.
- C/C++ remain without middle end optimizer.
- There is a necessity to design a common IR to share analysis and transformations across C/C++/Fortran.

Why CIL?

- LLVM is too low level: C++ STL constructs needs higher level IR to optimize effectively.
- Loop Transforms and Multi-dimensional arrays needs to analyzed accurately.
- Lack of optimizations for Parallel / Heterogeneous Programming constructs like OpenMP, OpenACC, etc.
- Need ABI independent IR

Limitations of LLVM IR (1) :C++ library optimizations

Found in hot part of a SPEC CPU 2017 benchmark omnetpp:

```
#include <map>

int func(std::map<int,int> &some_map) {

    for (auto &pair : some_map) {

    }

    return 0;

}
```

LLVM doesn't remove the dead loop

Limitations of LLVM IR (1) :C++ library optimizations

```
#include <string>

int main() {
    std::string str = "";
    if (str == "") {
        return 0;
    }
    return 1;
}
```

Limitations of LLVM IR (2) : Multi-dimensional arrays

- Linearized access using GEP
- Dilinearization is non-trivial effort
- Overcome by various methods like intrinsics, metadata, etc.
- Existing solutions doesn't fit well with LLVM IR.
- Leads to poor Dependence Analysis and hence Loop optimizations.

Limitations of LLVM IR (3) : Heterogeneous computing

- Various Heterogeneous programming paradigms supported in Clang
 - OpenMP
 - OpenACC
 - OpenCL
- Not natively supported in LLVM IR
- Mostly handled using runtime calls and intrinsics.
- Cross module optimizations
- Increasing use of C / C++ in DL / RL Frameworks / libraries.

CIL - Definition

Types

- Very close correspondence with C types.
- **Target independent**
- **Scalars:** `cil.int`, `cil.char` `cil.float`, `cil.long`,...
- Support for access / type qualifiers
- **Pointers:** `cil.pointer<cil.int>`, ...
- **Arrays:**
 - `!cil.array<10 x !cil.int>`
 - `!cil.array<10 x !cil.array<10 x !cil.int>>`
- **Derived Types:**
 - `!cil.struct.example<!cil.int, !cil.float>>`

Operations:

- Close correspondence with LLVM instructions with some exceptions:
 - Alloca and dealloca op
 - LoadOp / StoreOp
 - CILConstant
 - CILGlobalOp
 - **CILCastToMemRef**
 - **CILIfOp** (terminator operation)
- **Struct element access:**
 - `cil.struct_element %<struct_pointer>, %<field_index>`
- **Array element access:**
 - `cil.array_index %3, %4 : (!cil.pointer<!cil.array<10 x !cil.int>>, !cil.int) -> !cil.pointer<!cil.int>`
- **Pointer index access** (unbounded access via pointer type)
 - `cil.pointer_index %2, %1 : (!cil.pointer<!cil.int>, !cil.int) -> !cil.pointer<!cil.int>`

CIL Example 1:

```
struct example {  
    int a;  
    int b;  
};  
struct example e2;  
int a[10];  
int sum() { return a[3] + e2.a; }
```

```
module {  
    %0 = cil.global @e2 : !cil.struct.example<!cil.int, !cil.int>  
    %1 = cil.global @a : !cil.array<10 x !cil.int>  
    func @sum() -> !cil.int {  
        %2 = cil.global_address_of @a  
        %4 = cil.constant( 3 : i32 ): !cil.int  
        %5 = cil.array_index %2, %4  
        %6 = cil.load %5 : !cil.pointer<!cil.int> -> !cil.int  
        %7 = cil.global_address_of @e2  
        %8 = cil.global_address_of @e2  
        %9 = cil.constant( 0 : i32 ): !cil.int  
        %10 = cil.struct_element %8, %9  
        %11 = cil.load %10 : !cil.pointer<!cil.int> -> !cil.int  
        %12 = cil.addi %6, %11 : !cil.int  
        return %12 : !cil.int  
    }  
}
```

Loops

```
int a[100][100];
int b[100][100],c[100][100],i,j;
...
for(i=0; i<100; i++) {
    for(j=0; j<100; j++) {
        c[i][j] = a[i][j] + b[i][j];
    }
}
```

```
%2 = cil.alloca !cil.array<100 x !cil.array<100 x !cil.int>>
%3 = cil.alloca !cil.array<100 x !cil.array<100 x !cil.int>>
%4 = cil.alloca !cil.array<100 x !cil.array<100 x !cil.int>>
...
cil.for_loop %arg0 = %10, %11, %12 {
    %17 = cil.constant( 0 : i32 ): !cil.int
    %18 = cil.constant( 100 : i32 ): !cil.int
    %19 = cil.constant( 1 : i32 ): i32
    cil.for_loop %arg1 = %17, %18, %19 {
        %20 = cil.array_index %4, %arg0 :
        %21 = cil.array_index %20, %arg1 : !cil.pointer<!cil.int>
        %22 = cil.array_index %2, %arg0
        %23 = cil.array_index %22, %arg1 : !cil.pointer<!cil.int>
        %24 = cil.load %23 : !cil.int
        %25 = cil.array_index %3, %arg0
        %26 = cil.array_index %25, %arg1 : !cil.pointer<!cil.int>
        %27 = cil.load %26 : !cil.int
        %28 = cil.addi %24, %27 : !cil.int
        cil.store %28, %21 : !cil.pointer<!cil.int>
        %29 = cil.load %21 : !cil.int
    }
}
```

Unstructured Control Flow

- Unstructured control flow: break , continue, goto, ..
- Nested regions would cause problems
- Needs an analysis to convert them to Loop operations like loop.for, affine.for, etc.

C++ representation in CIL

Classes in CIL

```
class A {  
    public:  
    int a;  
    int b;  
    void func() { b = a + 3; }  
};
```

```
cil.class A {  
    cil.field_decl a : !cil.int  
    cil.field_decl b : !cil.int  
    func @_ZN1A4funcEv() attributes {original_name="func"} {  
        %2 = cil.cxx_this : !cil.pointer<!cil.class<A>>  
        %3 = cil.field_access %2, @b::@A  
        %4 = cil.cxx_this : !cil.pointer<!cil.class<A>>  
        %5 = cil.field_access %4, @a::@A  
        %6 = cil.load %5 : !cil.int  
        %7 = cil.constant( 3 : i32 ): !cil.int  
        %8 = cil.addi %6, %7 : !cil.int  
        cil.store %8, %3 : !cil.pointer<!cil.int>  
        %9 = cil.load %3 : !cil.int  
        return  
    }  
}
```


Lowering CIL Classes to Low Level CIL

- Converting high level CIL to low level CIL

```
cil.class A {
  cil.field_decl a : !cil.int
  cil.field_decl b : !cil.int
  func @_ZN1A4funcEv() {
    %2 = cil.cxx_this
    %3 = cil.field_access %2, @b::@A
    %4 = cil.cxx_this
    %5 = cil.field_access %4, @a::@A
    %6 = cil.load %5 : !cil.int
    %7 = cil.constant( 3 : i32 ): !cil.int
    %8 = cil.addi %6, %7 : !cil.int
    cil.store %8, %3 : !cil.pointer<!cil.int>
    %9 = cil.load %3 : !cil.int
    return
  }
}
```

```
func @_ZN1A4funcEv(%arg0:
!cil.pointer<!cil.struct.A<!cil.int, !cil.int>>) attributes
{original_name = "func"} {
  %2 = cil.constant( 1 : i32 ): !cil.int
  %3 = struct_element %arg0, %2 : cil.pointer<!cil.int>
  %4 = cil.constant( 0 : i32 ): !cil.int
  %5 = struct_element %arg0, %4 : !cil.pointer<!cil.int>
  %6 = cil.load %5 : !cil.int
  %7 = cil.constant( 3 : i32 ): !cil.int
  %8 = cil.addi %6, %7 : !cil.int
  cil.store %8, %3 : !cil.pointer<!cil.int>
  %9 = cil.load %3 : !cil.int
  return
}
```

Lowering CIL Classes to LLVM

- Converting low level CIL to LLVM

```
func @_ZN1A4funcEv(%arg0:
!cil.pointer<!cil.struct.A<!cil.int, !cil.int>>)
attributes {original_name = "func"} {
    %2 = cil.constant( 1 : i32 ): !cil.int
    %3 = struct_element %arg0, %2
    %4 = cil.constant( 0 : i32 ): !cil.int
    %5 = struct_element %arg0, %4
    %6 = cil.load %5 : !cil.int
    %7 = cil.constant( 3 : i32 ): !cil.int
    %8 = cil.addi %6, %7 : !cil.int
    cil.store %8, %3 : !cil.pointer<!cil.int>
    %9 = cil.load %3 : !cil.int
    return
}
```

```
%0 = type { i32, i32 }

; Function Attrs: nofree norecurse nounwind
define void @_ZN1A4funcEv(%0* nocapture %0)
local_unnamed_addr #0 !dbg !3 {
    %2 = getelementptr %0, %0* %0, i64 0, i32 1, !dbg !7
    %3 = getelementptr %0, %0* %0, i64 0, i32 0, !dbg !9
    %4 = load i32, i32* %3, align 4, !dbg !9
    %5 = add i32 %4, 3, !dbg !10
    store i32 %5, i32* %2, align 4, !dbg !10
    ret void, !dbg !11
}
```

Templates in CIL

```
template <class T> class A {  
public:  
    T a;  
};
```

```
A<int> o;  
A<float> o2;
```

```
cil.class A.0 {  
    cil.field_decl a : !cil.float  
}  
cil.class A {  
    cil.field_decl a : !cil.int  
}
```

```
%2 = cil.alloca !cil.class<A> : !cil.pointer<!cil.class<A>>  
%3 = cil.alloca !cil.class<A.0> : !cil.pointer<!cil.class<A.0>>
```

Operator Overloading in CIL

```
class cout {
public:
    int dummy;
};

void
operator<<(cout &obj, int val)
{
    printf("%d \n", val);
}
```

```
cil.class cout {
    cil.field_decl dummy : !cil.int
}

func @_ZlsR4couti(%arg0:
!cil.ref_pointer<!cil.class<cout>>, %arg1: !cil.int) {
    %1 = cil.alloca !cil.ref_pointer<!cil.class<cout>>
    cil.store %arg0, %1
    %2 = cil.alloca !cil.int : !cil.pointer<!cil.int>
    cil.store %arg1, %2 : !cil.pointer<!cil.int>
    %3 = cil.global_address_of @__str_tmp0
    %4 = cil.pointer_bitcast %3
    %5 = cil.load %2 : !cil.int
    %6 = cil.call @printf(%4, %5) : !cil.int,
    return
}
```

Inheritance in CIL

```
class C {
private:
    int c;

public:
    void set_c(int val) { c = val; }
    int get_c() { return c; }
};

class B : public C {
public:
    int e;
};
```

```
cil.class C {
    cil.field_decl c : !cil.int
    func @_ZN1C5set_cEi(%arg0: !cil.int) {
        %1 = cil.alloca !cil.int : !cil.pointer<!cil.int>
        cil.store %arg0, %1 : !cil.pointer<!cil.int>
        %2 = cil.cxx_this : !cil.pointer<!cil.class<C>>
        %3 = cil.field_access %2, @c::@C
        %4 = cil.load %1 : !cil.int
        cil.store %4, %3 : !cil.pointer<!cil.int>
        %5 = cil.load %3 : !cil.int
        return
    }
    func @_ZN1C5get_cEv() -> !cil.int {
        %1 = cil.cxx_this : !cil.pointer<!cil.class<C>>
        %2 = cil.field_access %1, @c::@C
        %3 = cil.load %2 : !cil.int
        return %3 : !cil.int
        ^bb1: // no predecessors
        cil.unreachable
    }
}

cil.class B inherits [!cil.class<C>] {
    cil.field_decl e : !cil.int
}
```

STL Optimizations using CIL (1)

```
void func() {  
    std::vector<int> vec;  
    vec.push_back(1);  
    vec.push_back(2);  
    vec.push_back(3);  
  
    vec.insert(vec.end(), {4, 5, 6});  
}
```

```
void func() {  
    std::vector<int> vec;  
    vec.insert(vec.end(), {1, 2, 3});  
    vec.insert(vec.end(), {4, 5, 6});  
}
```

```
%15 = cil.alloc @cil.class<std::__1::vector>,
@_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector() :
!cil.pointer<!cil.class<std::__1::vector>>
%16 = cil.constant( 1 : i32 ): !cil.int
cil.member_call %15, @_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector(%16)
%17 = cil.constant( 2 : i32 ): !cil.int
cil.member_call %15, @_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector(%17)
%18 = cil.constant( 3 : i32 ): !cil.int
cil.member_call %15, @_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector(%18)
%19 = cil.alloc @cil.class<std::__1::__wrap_iter> :
!cil.pointer<!cil.class<std::__1::__wrap_iter>>
%20 = cil.load %19 : !cil.class<std::__1::__wrap_iter>
%21 = cil.global { [4 : i32, 5 : i32, 6 : i32] } : !cil.class<std::initializer_list>
{sym_name = "", value = [4 : i32, 5 : i32, 6 : i32]}
```

```
%15 = cil.alloc @cil.class<std::__1::vector>, @_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector() :
!cil.pointer<!cil.class<std::__1::vector>>
%16 = cil.global { [1 : i32, 2 : i32, 3 : i32] } : !cil.class<std::initializer_list> {sym_name = "", value = [1 :
i32, 2 : i32, 3 : i32]}
%17 = cil.alloc @cil.class<std::__1::__wrap_iter> : !cil.pointer<!cil.class<std::__1::__wrap_iter>>
%18 = cil.load %17 : !cil.class<std::__1::__wrap_iter>
%19 = cil.member_call %15,
@_ZNSt3__1vectorIiNS_9allocatorIiEEEC1Ev::@vector(%18, %16) :
!cil.class<std::__1::__wrap_iter>,
%20 = cil.alloc @cil.class<std::__1::__wrap_iter> : !cil.pointer<!cil.class<std::__1::__wrap_iter>>
%21 = cil.load %20 : !cil.class<std::__1::__wrap_iter>
%22 = cil.global { [4 : i32, 5 : i32, 6 : i32] } : !cil.class<std::initializer_list> {sym_name = "", value = [4
: i32, 5 : i32, 6 : i32]}
```

STL Optimizations using CIL (2)

- Adding `std::vector::reserve()` before loop

```
void func() {  
    std::vector<int> vecA;  
  
    for(int i=0; i<10; i++) {  
        vecA.push_back(i);  
    }  
}
```

```
void func() {  
    std::vector<int> vecA;  
  
    vecA.reserve(10);  
    for(int i=0; i<10; i++) {  
        vecA.push_back(i);  
    }  
}
```



```
^bb2: // 2 preds: ^bb1, ^bb4
    %11 = cil.load %9 : !cil.int
    %12 = cil.constant( 10 : i32 ): !cil.int
    %13 = cil.cmpi slt %11, %12 : !cil.bool
    cil.for %13, ^bb3, ^bb5
^bb3: // pred: ^bb2
    cil.member_call %7,
@_ZNSt3__16vectorIiNS_9allocatorIiEEE9push_backERKi::@vector(%9)
    br ^bb4
^bb4: // pred: ^bb3
    %14 = cil.load %9 : !cil.int
    %15 = cil.constant( 1 : i32 ): !cil.int
    %16 = cil.addi %14, %15 : !cil.int
    cil.store %16, %9 : !cil.pointer<!cil.int>
    br ^bb2
```

```
    %9 = cil.alloca !cil.int :
!cil.pointer<!cil.int>
    %10 = cil.constant( 0 : i32 ): !cil.int
    cil.store %10, %9 : !cil.pointer<!cil.int>
    br ^bb1
^bb1: // 2 preds: ^bb0, ^bb2
    %11 = cil.constant( 10 : i32 ): !cil.int
    cil.member_call %7,
@_ZNSt3__16vectorIiNS_9allocatorIiEEE7reserveEm::@
vector(%11)
    %12 = cil.load %9 : !cil.int
    %13 = cil.constant( 10 : i32 ): !cil.int
    %14 = cil.cmpi slt %12, %13 : !cil.bool
    cil.for %14, ^bb2, ^bb3
^bb2: // pred: ^bb1
    cil.member_call %7,
@_ZNSt3__16vectorIiNS_9allocatorIiEEE9push_backERKi::@vector(%9)
    %15 = cil.load %9 : !cil.int
    %16 = cil.constant( 1 : i32 ): !cil.int
    %17 = cil.addi %15, %16 : !cil.int
    cil.store %17, %9 : !cil.pointer<!cil.int>
```

Current Progress in C++

- Major support features include:
 - **C part of C++**
 - **Templates**
 - **Inheritance**
 - **Overloading**
 - **Namespaces.**
- Considerable progress in compiling STL.
 - **Partial support for `iostream`, `vector`, `string`, `set` ,...**
- Compiles and runs 80+ cpp unit tests.
- Compiling and Running a SPEC CPU 2017 C++ benchmark

CIL and Fortran IR

CIL and Fortran

- FC - Fortran compiler, now emits CIL from high level fortran dialect.
- Multiple levels of IR.
- Optimisations at different abstractions.
- CIL is further lowered to LLVM dialect.
- All CIL optimisations can be reused
- FC emits CIL for 2 SPEC CPU 2017 benchmarks.
- Experimental LTO framework.

FC Fortran IR Example

```
program hello
  print *, "hello"
end program
```



```
module {
  fc.function @hello() -> !cil.int {
    %0 = fc.constant_string("hello^@") : !fc.array<0:5 *
!cil.char8>
    fc.print %0:!fc.array<0:5 * !cil.char8>
    %1 = cil.constant( 0 : i32 ): !cil.int
    fc.return %1
  }
}
```

Fortran IR: Nested functions

```
subroutine sub1
  integer :: a
  print *, func()
  contains
  integer function func
    func = a
  end function func
end
```



```
fc.function @sub1() {
  fc.function @func() -> i32 {
    %2 = fc.get_element_ref @a::@sub1 : !fc.ref<i32>
    %3 = fc.allocate func : !fc.ref<i32>
    %4 = fc.load %2 {name = "a"} : i32
    fc.store %4, %3 {name = "func"} : !fc.ref<i32>
    %5 = fc.load %3 : i32
    fc.return %5
  }
  %0 = fc.allocate a, implicitly_captured !fc.ref<i32>
  %1 = fc.call @func::@sub1() : i32,
  fc.print %1 {arg_info = #fc.is_string< >}
  fc.return
}
```

Fortran CIL Example

```
program hello
  integer :: a
  a = 10
  print *, a + 15
end program
```



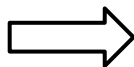
```
module {
  %0 = cil.global @fc_internal_argv {sym_name =
"fc_internal_argv"} : !cil.pointer<!cil.pointer<!cil.char8>>
  %1 = cil.global @fc_internal_argc : !cil.int
  func @hello() -> !cil.int {
    %2 = cil.alloca !cil.int : !cil.pointer<!cil.int>
    %3 = cil.constant( 10 : i32 ): !cil.int
    cil.store %3, %2 : !cil.pointer<!cil.int>
    %4 = cil.load %2 : !cil.pointer<!cil.int> -> !cil.int
    %5 = cil.constant( 15 : i32 ): !cil.int
    %6 = cil.addi %4, %5 : !cil.int
    %7 = cil.constant( 3 : i32 ): !cil.int
    %8 = cil.constant( 2 : i32 ): !cil.int
    cil.call @_fc_runtime_print(%8, %7, %6)
    %9 = cil.constant( 0 : i32 ): !cil.int
    cil.return %9 : !cil.int
  }
  ...
}
```

CIL - LTO

- Link multiple CIL modules into single CIL module.
- Resolves functions declarations with their definitions.
- Resolves global variable declarations.
- LTO helps to implements IPOs such as inliner, data layout optimisations in CIL

CIL LTO Example

```
module {
  %0 = cil.global @__str_tmp0 {constant,
    value = " %d \0A\00"} : ...
  func @main() -> !cil.int {
    ...
  }
  ...
  func @add(!cil.int, !cil.int) -> !cil.int
}
```



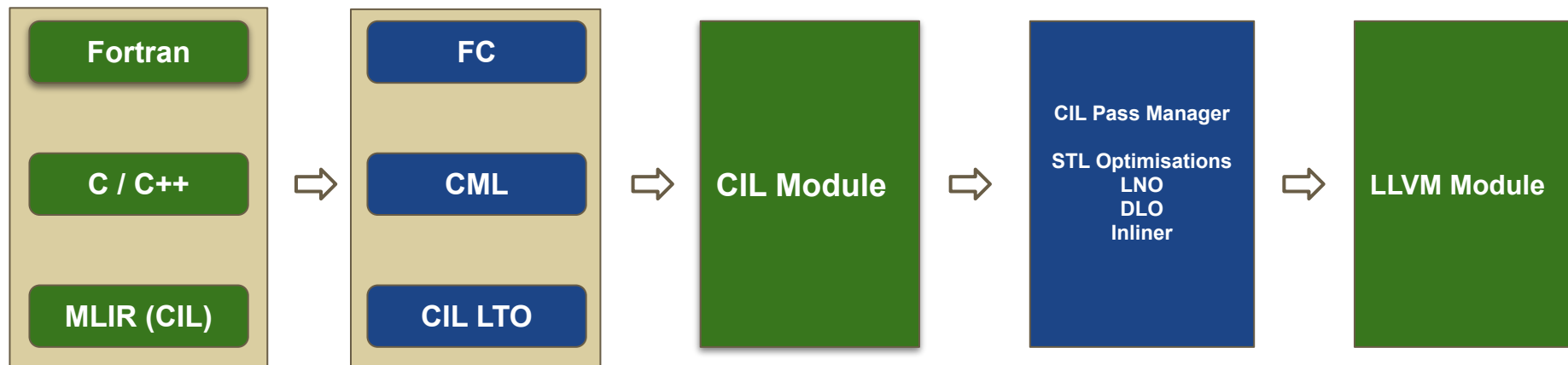
```
module {
  %0 = cil.global @__str_tmp0 {constant, value = " %d
\0A\00"} : ...
  func @printf(!cil.pointer<!cil.char>) -> !cil.int
  func @main() -> !cil.int {
    ...
  }
  func @add(%arg0: !cil.int, %arg1: !cil.int) ->
!cil.int {
    ...
  }
}
```

```
module {
  func @add(%arg0: !cil.int, %arg1: !cil.int)
-> !cil.int {
    ...
  }
}
```

Data Layout Optimisations in CIL

- Instance interleaving and Dead field elimination optimisations are implemented in CIL
- Runs as LTO pass.
- Identification of struct access is simpler as compared to LLVM because there is separate operation for struct access.
- Approximately 35% improvement is seen in one of SPEC CPU 2017 benchmark.

C, C++, Fortran and CIL



CIL - Applications

- Optimising source codes for C, C++ and Fortran
- Optimising tensorflow graphs
- Optimising ONNX models
- Custom Hardware

Current Status

- Explored
 - C99
 - C++ : templates, operator overloading, Inheritance, ..
 - C++ STL
 - Fortran: low level ABI agnostic IR with LTO
 - Loop / Data layout optimization capability
- Yet to see:
 - Clang Integration
 - Clang static analysis toolchain
 - C++-11 and later standards
 - Other languages in Clang
 - C++ exceptions

Roadmap

CIL basic dialect to be open sourced soon

More focus on C++

Vectorizer integrated to Loop Nest Optimizer is work in progress.

LTO integration with Clang driver

Thank you
